

$p$  prime

$$1^2, 2^2, 3^2, \dots, \left(\frac{p-1}{2}\right)^2$$

are all distinct mod  $p$ .

Because if  $i^2 \equiv j^2 \pmod{p}$

then  $i^2 - j^2$  is divisible by  $p$

and  $(i-j)(i+j)$

if  $i, j$  are both  $\leq \frac{p-1}{2}$

then  $i-j$  &  $i+j$  are  $< p$

$\Rightarrow i-j = 0$  because  $p$

does not divide  $i+j$  so it must divide  $i-j$  and since

the only  $\# < p$  that  $p$  divides is 0.

Recall  $i^2 \equiv (p-i)^2 \equiv p^2 - 2pi + i^2 \equiv$

So all the quadratic residues are  
 $\{1^2, 2^2, 3^2, \dots, \left(\frac{p-1}{2}\right)^2\}$

$$\begin{array}{r} 46 \\ 17 \overline{) 787} \\ \underline{782} \\ 5 \end{array}$$

$$\begin{array}{r} 46 \\ 17 \\ \underline{322} \\ 46 \end{array}$$

$$J(17, 787) = J(787 \bmod 17, 17) \cdot (-1)^{\frac{16 \cdot 786}{4}}$$

$$= J(5, 17) \cdot 1$$

$$= J(17 \bmod 5, 5) \cdot (-1)^{\frac{4 \cdot 16}{4}}$$

$$= J(2, 5)$$

$$= J\left(\frac{2}{2}, 5\right) \cdot (-1)^{\frac{(2 \cdot 5 - 1)}{2}}$$

$$= -J(1, 5) = -1$$

$$17^{393} \pmod{787} \equiv -1$$

$$5^8 \equiv -1 \pmod{17}$$

$$17^{395} = 17^{393} \cdot 17^2 \equiv -17^2 \pmod{787}$$

## Jacobi Symbol

there is  
no solution  
to  $17 \equiv x^2 \pmod{787}$

$$17 \pmod{787} = 17^{393} \pmod{787} = 17^{\frac{787-1}{2}} \pmod{787} = \left(\frac{17}{787}\right) = -1$$

We start with the Legendre symbol

$$\left(\frac{a}{p}\right) = \begin{cases} 1 & \text{if } a \in QR[p] \\ -1 & \text{if } a \notin QR[p] \end{cases}$$

and for

$$n = p_1 p_2 \cdots p_k$$

we set

$$J(a, n) = \left(\frac{a}{p_1}\right) \left(\frac{a}{p_2}\right) \cdots \left(\frac{a}{p_k}\right)$$

However, for  $n$  odd, we have

$$J(a, n) = \begin{cases} 1 & \text{if } a = 1 \\ J(a/2, n)(-1)^{(n^2-1)/8} & \text{if } a \text{ is even} \\ J(n \bmod a, a)(-1)^{(n-1)(a-1)/4} & \text{if } a > 1 \text{ and odd} \end{cases}$$

Take a bunch of samples

$$a_1, a_2, \dots, a_k$$

and test

$$J(a_1, n) \quad \text{vs.} \quad a_1^{\frac{n-1}{2}} \pmod{n}$$

$$J(a_2, n) \quad \text{vs.} \quad a_2^{\frac{n-1}{2}} \pmod{n}$$

$$J(a_3, n) \quad \text{vs.} \quad a_3^{\frac{n-1}{2}} \pmod{n}$$

⋮

$$J(a_k, n) \quad \text{vs.} \quad a_k^{\frac{n-1}{2}} \pmod{n}$$

If  $a_1, \dots, a_k$  were chosen at random then because ~~the~~  $J(a_i, n) \not\equiv a_i^{\frac{n-1}{2}} \pmod{n}$  for half of the integers  $a$ , there is a  $(\frac{1}{2})^k$  chance that my  $a_i$ 's were chosen so that  $J(a_i, n) \equiv a_i^{\frac{n-1}{2}} \pmod{n}$  even though  $n$  is not prime.

## Quadratic Sieve

**Example:**  $m = 91$

$$23^2 \equiv 16^2 \pmod{91}$$

$$(23+16) = 39 \quad 23-16 = 7$$

$a$	19	1	23	18	2	24	16
$a^2$	88	1	74	51	4	30	74

$$\begin{aligned} 91 &= \gcd(91, 23 + 16) \times \gcd(91, 23 - 16) \\ &= \gcd(91, 39) \times \gcd(91, 7) \\ &= 13 \times 7 \end{aligned}$$

## Exercises

1. An individual publishes an RSA modulus of  $m = 350123$  and an encryption exponent  $e = 37$ . Find his decrypting exponent, given that one of the factors of  $m$  is 347.
2. Encrypt each letter of the word **BANG** individually using the RSA system with  $m = 143$  and  $e = 7$ . In translating letters into numbers, send **A** to 10, **B** to 11, ..., **Z** to 35.
3. Using the same system described in the previous problem, find the decrypting exponent  $d$  and decode the message 132 (a single letter).
4. Factor  $m = 773,771$  into the product of two primes given that  $\phi(m) = 771,552$ .

$$\phi(m) = (p-1)(q-1) = pq - p - q + 1$$

$\begin{matrix} \text{771552} & & & & \text{= 773771} \\ \parallel & & & & \\ \text{771552} & & & & \end{matrix}$

$$p + q = 773771 - 771552 + 1$$

$$p + q = 2220$$

$$p + \frac{773771}{p} = 2220$$

$$p = \frac{2220 \pm \sqrt{2220^2 - 4 \cdot 773771}}{2}$$

$$= 1787, 433$$

$$p^2 - 2220p + 773771 = 0$$

# Diffie-Hellman Public Key Exchange

1. People  $P_1, P_2, \dots, P_k$  agree on a modulus  $p$  in which they agree to do their calculations.
2. They also agree on a common base,  $a$ , which must be a primitive root of  $p$
3. Each person  $P_i$  secretly selects a number,  $S_i$ , from 1 to  $p - 1$  and publicly announces the value  $\beta_i = a^{S_i} \bmod p$ .

Alice  $a^{S_1} \longrightarrow E_1(x)$   
Bob  $a^{S_2} \longrightarrow E_2(x)$   
they use common key  $a^{S_1 S_2} = (a^{S_1})^{S_2} = (a^{S_2})^{S_1}$