

Basics of tableaux

July 24, 2003

Contents

1	Permutations	1
2	Partitions	2
2.1	Ferrers Diagrams	3
3	Tableaux	3
3.1	Counting standard tableaux	4
4	Robinson-Schensted Correspondence	4
4.1	Insertion/deletion algorithm	4
4.2	The correspondence	5
4.3	Direct consequences	5
5	Plactic monoid	6
5.1	Knuth equivalence	6
5.2	Knuth slow insertion	7
6	The tableaux monoid and the plactic algebra	8
6.1	Multiplication in the monoid	8
6.2	A result on bumping	9
6.3	The plactic algebra and Pieri rule	9

1 Permutations

A **permutation** of the integers $\{1, \dots, n\}$ is a one-to-one map of the integers onto themselves: $i \rightarrow \sigma_i$. There are several ways to denote such an action.

Two line notation: $\sigma = \begin{pmatrix} 1 & 2 & 3 & \cdots & n \\ \sigma_1 & \sigma_2 & \sigma_3 & \cdots & \sigma_n \end{pmatrix}$ **One line notation:** $\sigma = (\sigma_1 \sigma_2 \dots \sigma_n)$

The **symmetric group** S_n is the set of permutations of the integers $\{1, \dots, n\}$ under composition of maps. ListPerm(3);

$$S_3 = \{(123), (132), (312), (213), (231), (321)\}$$

MultPerm([2,1,3],[3,1,2]);

$$\sigma = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} \text{ and } \tau = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix} \implies \tau\sigma = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 3 & 2 \end{pmatrix}$$

The symmetric group is not abelian in general. That is, the elements do not necessarily commute.

$$\sigma = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} \text{ and } \tau = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix} \implies \sigma\tau = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix}$$

The **inverse** of $\sigma \in S_n$ is the element σ^{-1} where $\sigma^{-1}\sigma = id$.

$$\sigma = \begin{pmatrix} 1 & 2 & 3 & \cdots & n \\ \sigma_1 & \sigma_2 & \sigma_3 & \cdots & \sigma_n \end{pmatrix}$$

$$\sigma^{-1} = \begin{pmatrix} \sigma_1 & \sigma_2 & \sigma_3 & \cdots & \sigma_n \\ 1 & 2 & 3 & \cdots & n \end{pmatrix} = \begin{pmatrix} 1 & 2 & 3 & \cdots & n \\ \sigma'_1 & \sigma'_2 & \sigma'_3 & \cdots & \sigma'_n \end{pmatrix}$$

The inverse of $\sigma = (32514)$ is

$$\sigma^{-1} = \begin{pmatrix} 3 & 2 & 5 & 1 & 4 \\ 1 & 2 & 3 & 4 & 5 \end{pmatrix} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 4 & 2 & 1 & 5 & 3 \end{pmatrix}$$

A permutation that is its own inverse is called an **involution**. Note that $\sigma = \sigma^{-1} \iff \sigma\sigma = id$. The involutions of S_3 are (123) , (213) , (321) , (132) . For example, check (213) :

$$\begin{pmatrix} 2 & 1 & 3 \\ 1 & 2 & 3 \end{pmatrix} \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} = \begin{pmatrix} 2 & 1 & 3 \\ 1 & 2 & 3 \end{pmatrix} (2,1,3) = (1,2,3)$$

There is a third notation used in the study of permutations called **cycle notation**. $\sigma = (a_1, \dots, a_\ell) \cdots (b_1, \dots, b_j)$ means $\sigma(a_i) = a_{i+1}$

$$\sigma = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} = (213) = (1,2)(3)$$

The **cycle structure** of a permutation $\sigma \in S_n$ is the non-increasing vector of integers determined by the number of elements in the disjoint cycles of σ .

2 Partitions

The conjugacy classes of permutations are closely related to their cycle structure. This is one of many reasons to closely study such vectors. A **partition** is a vector of integers $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_n)$ where $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n > 0$. The partitions of 4 are

$$(4), (3,1), (2,2), (2,1,1), (1,1,1,1)$$

Alternative Notation: $\lambda = (\lambda_1^{m_1}, \lambda_2^{m_2}, \dots, \lambda_\ell^{m_\ell})$

λ_i occurs with multiplicity m_i

there are m_i occurrences of λ_i .

In this case, the partitions of 4 are denoted

$$(4), (3,1), (2^2), (2,1^2), (1^4)$$

The **degree** or **order** of a partition λ is $|\lambda| = \lambda_1 + \lambda_2 + \dots + \lambda_n$. We say “ λ is a partition of $|\lambda|$ ” and denote this by $\lambda \vdash |\lambda|$. $\lambda \vdash 4$ is to say $|\lambda| = 4$ and $\lambda = (2,1,1)$ is a partition of 4. The **length** of a partition is the number of parts. The length of λ is denoted $\ell(\lambda)$. $\lambda = (2,1,1) \implies \ell(\lambda) = 3$. The **conjugate** of a partition λ is $\lambda' = (\lambda'_1, \lambda'_2, \dots, \lambda'_n)$ where $\lambda'_i =$ number of parts of $\lambda \geq i$. 3 parts of $(2,1,1)$ are ≥ 1 implies $\lambda'_1 = 3$, 1 part of $(2,1,1)$ is ≥ 2 implies $\lambda'_2 = 1$. Thus $\lambda' = (3,1)$.

2.1 Ferrers Diagrams

The **Ferrers diagram** of a partition λ is the collection of stacked boxes, arranged in left-justified rows, so that number of boxes in each row is λ_i .

$$\lambda = (5, 2, 2, 1) = \begin{array}{cccc} \square & & & \\ \square & \square & & \\ \square & \square & & \\ \square & \square & \square & \square \end{array}$$

Everything translates into the notation of shapes. The total number of boxes gives the degree of the partition. The partition in the previous example has degree 10. The number of rows is the length of the partition, $\ell(\lambda) = 4$. The reflection of the diagram about $y = x$ is the conjugate partition

$$\lambda' = \begin{array}{cccc} \square & & & \\ \square & \square & & \\ \square & \square & \square & \square \\ \square & \square & \square & \square \end{array} \implies \lambda' = (4, 3, 1, 1, 1) = (4, 3, 1^3).$$

The partitions of 4 are



3 Tableaux

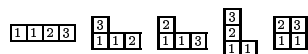
The symmetric group is linked to beautiful combinatorics. In particular, a combinatorial structure obtained by putting numbers into Ferrers diagrams is of utmost importance. A few examples of the role these objects (called tableaux) play are:

- There is a bijection between permutations and pairs of standard tableaux with the same shape (or more generally between two rowed arrays and pairs of tableaux).
- The number of involutions of $\{1, 2, \dots, n\}$ is the number of standard tableaux of n
- The tableaux form an associative monoid that facilitates the study of symmetric functions.

A **semi-standard tableau** of shape λ and content μ is a filling of the Ferrers diagram for λ with μ_1 ones, μ_2 twos, ... so that the numbers are weakly increasing in rows and strictly increasing in columns.



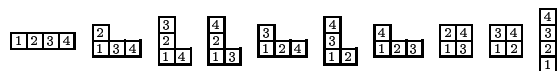
is a tableau of shape $\lambda = (5, 2, 2, 1)$ and content $(3, 3, 2, 1, 1)$. All the semi-standard tableaux of degree 4 with content $(2, 1, 1)$, that is with 2 ones, 1 two, 1 three are:



A **standard tableau** of shape λ is a semi-standard tableau with content $(1, 1, \dots, 1)$. In a standard tableau, the integers $\{1, 2, \dots, |\lambda|\}$ each occur once in the Ferrers diagram filled strictly increasing in rows and columns.



is a standard tableau of shape $\lambda = (5, 2, 2, 1)$ and degree 10. The set of all standard tableaux of degree 4 is

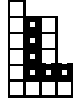


3.1 Counting standard tableaux

The number of standard tableaux of a given shape can be determined using the "hook formula". If $s = (i, j)$ is a cell in the diagram of λ , then it has a **hook**

$$H_s = H_{i,j} = \{(i, j') | j' \geq j\} \cup \{(i', j) | i' \geq i\}$$

with corresponding **hook-length** $h_s(\lambda) = |H_s|$. For example $h_{2,2}(4^2, 3^3, 1) = 6$ since $H_{2,2}$ are the outlined cells in



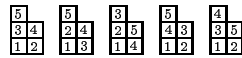
The hook formula: The number of standard tableaux of shape λ (f^λ) is $n!$ divided by the product of all hook-lengths;

$$f^\lambda = \frac{n!}{\prod_{s \in \lambda} h_s}$$

The number of standard tableaux of shape $(2, 2, 1)$ is determined by finding each hook-lengths h_s of λ (written in the corresponding cells s):

$$\begin{array}{|c|c|c|} \hline 1 & & \\ \hline 3 & 1 & \\ \hline 4 & 2 & \\ \hline \end{array} \implies f^{2,2,1} = \frac{5!}{4 \cdot 3 \cdot 2 \cdot 1 \cdot 1} = 5$$

The five standard tableaux of shape $(2, 2, 1)$ are

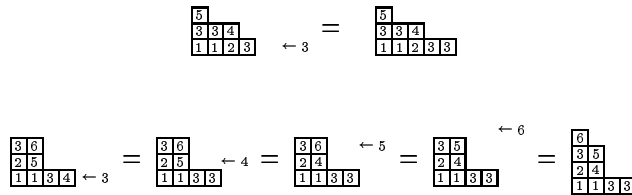


4 Robinson-Schensted Correspondence

4.1 Insertion/deletion algorithm

A fundamental operation on tableaux is given by the **Schensted insertion algorithm**, which sends a tableau T and a positive integer x to a new tableau that has one more box than T . This new tableau is denoted $T \leftarrow x$.

If x is weakly larger than every entry in the bottom row of T , add x to the end of this row. If not, find the leftmost entry e in the bottom row that is strictly larger than x . Replace this entry by x . Repeat the process on the next row with the letter e .



Note that this algorithm will send any word to a tableau by successively inserting the letters using Schensted's insertion algorithm.

$$\begin{aligned} \emptyset \leftarrow 234412232 &= \begin{array}{|c|} \hline 2 \\ \hline \end{array} \leftarrow 34412232 = \begin{array}{|c|c|} \hline 2 & 3 \\ \hline \end{array} \leftarrow 4412232 = \dots = \begin{array}{|c|c|c|c|} \hline 2 & 3 & 4 & 4 \\ \hline \end{array} \leftarrow 12232 \\ &= \begin{array}{|c|c|c|c|} \hline 2 & & & \\ \hline 1 & 3 & 4 & 4 \\ \hline \end{array} \leftarrow 12232 = \begin{array}{|c|c|} \hline 2 & 3 \\ \hline 1 & 2 & 4 & 4 \\ \hline \end{array} \leftarrow 232 = \begin{array}{|c|c|c|} \hline 2 & 3 & 4 \\ \hline 1 & 2 & 2 & 4 \\ \hline \end{array} \leftarrow 32 = \begin{array}{|c|c|c|c|} \hline 2 & 3 & 4 & 4 \\ \hline 1 & 2 & 2 & 3 \\ \hline \end{array} \leftarrow 2 = \begin{array}{|c|c|c|c|} \hline 4 & & & \\ \hline 2 & 3 & 3 & 4 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array} \end{aligned}$$

Given that $\hat{T} = T \leftarrow x$ has a new cell in position c , we can easily reverse the process of insertion starting from the entry in position c . The **deletion algorithm** will send a given tableau T and a corner c of T to a tableau of the shape of T with the cell c removed.

Delete the entry x from corner c . If corner c was in the bottom row of T , we are done.

If not, find the rightmost entry e smaller than x in the row below that with c and replace e by x . Repeat this process with entry e .

If c is the cell containing 6 in the tableau T :

$$T = \begin{array}{|c|c|c|} \hline 6 & & \\ \hline 3 & 5 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|} \hline 3 & 5 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \xleftarrow{6} \rightarrow \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \xleftarrow{5} \rightarrow \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 5 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \xleftarrow{4} \rightarrow \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 5 & \\ \hline 1 & 1 & 3 & 4 \\ \hline \end{array}$$

4.2 The correspondence

A bijection between permutations and pairs of standard tableaux with the same shape arises from the insertion algorithm. The **Robinson-Schensted algorithm** is a method for constructing a pair of tableaux of the same shape from a permutation.

From any $\sigma \in S_n$, the first tableau in our pair is constructed from σ using the Schensted insertion algorithm. However, during this construction when a new box is added by inserting σ_i , put the letter i into this box in the second tableau. The first tableau is called the **insertion tableau** and the second is the **recording tableau**.

To find the pair of tableaux associated to $\sigma = 415326$, start by row inserting $\sigma_1 = 4$ into \emptyset ;

$$\begin{array}{l} \emptyset \leftarrow 4 = \begin{array}{|c|} \hline 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 1 \\ \hline \end{array} \\ \\ \begin{array}{|c|} \hline 4 \\ \hline \end{array} \leftarrow 1 = \begin{array}{|c|} \hline 4 \\ \hline 1 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 2 \\ \hline 1 \\ \hline \end{array} \\ \\ \begin{array}{|c|} \hline 4 \\ \hline 1 \\ \hline \end{array} \leftarrow 5 = \begin{array}{|c|} \hline 4 \\ \hline 1 & 5 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 2 \\ \hline 1 & 3 \\ \hline \end{array} \\ \\ \begin{array}{|c|} \hline 4 \\ \hline 1 & 5 \\ \hline \end{array} \leftarrow 3 = \begin{array}{|c|} \hline 4 & 5 \\ \hline 1 & 3 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 2 & 4 \\ \hline 1 & 3 \\ \hline \end{array} \\ \\ \begin{array}{|c|} \hline 4 & 5 \\ \hline 1 & 3 \\ \hline \end{array} \leftarrow 2 = \begin{array}{|c|} \hline 4 \\ \hline 3 & 5 \\ \hline 1 & 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 5 \\ \hline 2 & 4 \\ \hline 1 & 3 \\ \hline \end{array} \\ \\ \begin{array}{|c|} \hline 4 \\ \hline 3 & 5 \\ \hline 1 & 2 \\ \hline \end{array} \leftarrow 6 = \begin{array}{|c|} \hline 4 \\ \hline 3 & 5 \\ \hline 1 & 2 & 6 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 5 \\ \hline 2 & 4 \\ \hline 1 & 3 & 6 \\ \hline \end{array} \end{array}$$

Schensted($w[4,1,5,3,2,6]$);

Note that the recording tableau is a standard tableau of n since Q is formed by adding elements in increasing order to cells on the periphery.

Since the Schensted insertion algorithm is reversible, this process can be inverted to obtain a permutation of S_n from a pair of standard tableaux of the same shape. That is, starting with the entry e in the insertion tableau that lies in position n determined by the recording tableaux, we use the deletion process to remove this letter. Then repeat starting with $n - 1$. Thus, the **Robinson-Schensted Correspondence** is a bijection between elements of S_n and the set of ordered pairs (P, Q) of standard tableaux of n having the same shape.

4.3 Direct consequences

Theorem 1. The number of permutations of $\{1, 2, \dots, n\}$ is equal to the number of pairs of standard tableaux of the same shape λ as λ varies over all partitions of n .

$$\sum_{\lambda \vdash n} (f^\lambda)^2 = n!$$

When we consider only permutation that are involutions, another beautiful identity holds. Notice the number of involutions of S_4 :

$$(1\ 2\ 3\ 4), (2\ 1\ 3\ 4), (3\ 2\ 1\ 4), (4\ 2\ 3\ 1), (1\ 3\ 2\ 4), (1\ 4\ 3\ 2), (1\ 2\ 4\ 3), (2\ 1\ 4\ 3), (3\ 4\ 1\ 2), (4\ 3\ 2\ 1)$$

equals the number of standard tableaux of degree 4:

$$\begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 4 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 2 \\ \hline 1 & 3 & 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 3 \\ \hline 2 \\ \hline 1 & 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 4 \\ \hline 2 \\ \hline 1 & 3 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 3 & & 4 \\ \hline 1 & 2 & 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 4 \\ \hline 3 \\ \hline 1 & 2 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 4 & \\ \hline 1 & 2 & 3 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 2 & 4 \\ \hline 1 & 3 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 3 & 4 \\ \hline 1 & 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 4 \\ \hline 3 \\ \hline 2 \\ \hline 1 \\ \hline \end{array}.$$

This follows from the property: if σ corresponds to the pair of standard tableaux (P, Q) under the RS-correspondence, then σ^{-1} corresponds to the pair (Q, P) ¹. Thus for $\sigma = \sigma^{-1}$ an involution, $(P, Q) = (Q, P)$. That is, σ corresponds to a pair of the same standard tableaux.

Theorem 2. *The number of involutions on $\{1, \dots, n\}$ is the number of standard tableaux on n letters.*

5 Plactic monoid

Another important idea linked closely to Schensted insertion is the study of tableaux as a monoid. Recall that an associative **monoid** is a set G containing an identity e ($x \cdot e = e \cdot x = x$ for all $x \in G$) and having a binary operation (called multiplication) for which the associative law holds. That is, a monoid is a group without an inverse.

The collection of all words on a given alphabet forms a monoid called the **free monoid**, where multiplication is defined by word juxtaposition:

$$w = (3416647) \quad \text{and} \quad w' = (2155) \quad \text{give} \quad ww' = (34166472155)$$

The empty word \emptyset is the identity element. In particular, we have the monoid F on letters $\{1, \dots, m\}$. Using the obvious structure of the free monoid, a connection between words and tableaux that will enable us to induce a monoid on tableaux.

5.1 Knuth equivalence

The **row word** of tableau T , denoted $w(T)$, is the sequence of letters read from a tableau like a book in English.

$$w\left(\begin{array}{|c|c|} \hline 3 & 5 \\ \hline 2 & 4 \\ \hline 1 & 1 & 2 & 2 & 6 \\ \hline \end{array}\right) = 35241126$$

Note that **every tableau** $T \rightarrow$ **a word** $w(T)$. The original tableau is recovered from its row word by noting that the rows end at each **descent** – any position in a word w where $w_i > w_{i+1}$:

$$35|24|1126 \quad \longrightarrow \quad \begin{array}{|c|c|} \hline 3 & 5 \\ \hline 2 & 4 \\ \hline 1 & 1 & 2 & 2 & 6 \\ \hline \end{array}$$

However, not every word comes from a tableaux. Only when a word that is cut at its descents has pieces of weakly increasing length, and when stacked have columns that are strictly increasing does the word form a tableau. Otherwise, we may find the following scenario:

$$357241126 \rightarrow 357|24|1126 \rightarrow \begin{array}{|c|c|c|} \hline 3 & 5 & 7 \\ \hline 2 & 4 & \\ \hline 1 & 1 & 2 & 2 & 6 \\ \hline \end{array}$$

Some words \rightarrow a tableau.

¹Appendix

To reconcile that there are “too many” words, we shall classify certain words as “the same”. The **Knuth relation A** is move on a three letter word:

$$y z x \Rightarrow y x z \iff \boxed{y \mid z} \leftarrow x = \begin{array}{|c|c|} \hline y \\ \hline x & z \\ \hline \end{array} \text{ if } x < y \leq z$$

The **Knuth relation B** is move on a three letter word:

$$x z y \Rightarrow z x y \iff \boxed{x \mid z} \leftarrow y = \begin{array}{|c|c|} \hline z \\ \hline x & y \\ \hline \end{array} \text{ if } x \leq y < z$$

An **elementary Knuth transformation** on a word applies one of A or B or their inverses to three successive letters of a word. Thus, words w and w' are **Knuth equivalent** if w can be obtained by applying a sequence of elementary Knuth transformations to w' .

The set $M = F/R$ where R is the equivalence relation generated by the Knuth relations (the Knuth equivalence classes of words) is closed under multiplication (juxtaposition). That is, $w, u \in M$ are such that $wu \equiv w'u'$ if $w \equiv w'$ and $u \equiv u'$. Thus the multiplication on F descends to a multiplication on the set M . This makes M into an associative monoid called the **plactic monoid**.

It will develop that **each equivalence class \rightarrow tableau**. Then, a bijection between equivalence classes and tableaux will reveal that tableaux form a monoid that is isomorphic to the plactic monoid.

5.2 Knuth slow insertion

Recall, the insertion algorithm produces a tableau $T \leftarrow x$ from a tableau T and letter x . This algorithm can instead be achieved systematically by applying a sequence of Knuth transformations A and B to the left of $w(T)x$.

Consider only the row word of the bottom row of T , $w = u_1 \cdots u_p x' v_1 \cdots v_q x$ where $x' > x \geq u_p$. Check that x is strictly smaller than its two preceding letters. If so, we move x forward by transposing with the entry to its right since $v_q \geq v_{q-1} > x$ implies the Knuth operation A moves x to the left of v_q .

$$122334\underline{2} \stackrel{A}{\equiv} 12233\underline{2}4 \stackrel{A}{\equiv} 1223\underline{2}34$$

This continues until x is no longer smaller, then x rests while x' is pushed forward by transposing with the entry to its left. That is, use transformation A until the configuration $u_1 \cdots u_p x' x v_1 \cdots v_q$ is reached. Then $u_p \leq x$ and $x' > x$ implies that operation B moves x' to the left of u_{p-1} .

$$122\underline{3}234 \stackrel{B}{\equiv} 12\underline{3}2234 \stackrel{B}{\equiv} 1\underline{3}22234 \stackrel{B}{\equiv} \underline{3}122234$$

This continues until x' has been moved to the beginning of the row word w of the first row of T . As such, the letter x has been Schensted inserted into the bottom row and x' has been bumped. This process is then repeated with x' at the end of the row word of the second row of T , amounting exactly to the insertion of x' into the second row, etc. Thus, the Schensted insertion of x into T can be achieved by Knuth transformations on $w(T)x$. That is,

Proposition 3. For any tableau T and positive integer x ,

$$w(T \leftarrow x) \equiv w(T)x$$

More generally, given any word $w = x_1 \cdots x_\ell$ and starting from the empty tableau $T = \emptyset$ we have $w(\emptyset \leftarrow x_1) \equiv x_1$, then with $T = \emptyset \leftarrow x_1$, $w(T \leftarrow x_2) \equiv x_1 x_2$, and so. Thus proving:

Proposition 4. Every word is Knuth equivalent to the word of a tableau. In particular,

$$x_1 x_2 \cdots x_\ell \equiv w(\emptyset \leftarrow x_1 \leftarrow \cdots \leftarrow x_\ell)$$

implying that the word $w = x_1 \cdots x_\ell$ is Knuth equivalent to the tableau $T = \emptyset \leftarrow w$.

In fact a stronger result holds². ...

Theorem 5. *Every word is Knuth equivalent to the word of a unique tableau. That is, each equivalence class contains exactly one row word of a tableau.*

All words that are Knuth equivalent can be associated to the same tableau and no others:

$$\begin{aligned} 235512232 &\equiv 235512322 \equiv 235513222 \equiv 235531222 \\ &\equiv 235351222 \equiv 253351222 \equiv 523351222 \equiv w\left(\begin{array}{|c|c|c|c|} \hline 5 & & & \\ \hline 2 & 3 & 3 & 5 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array}\right) \end{aligned}$$

FreePlaxClass(w[5,2,3,3,5,1,2,2]); Gives all elements of that are Knuth equivalent to 523351222.

6 The tableaux monoid and the plactic algebra

We have now seen that the map w from a tableau to the row word of this tableau is a bijection between the set of tableaux and the Knuth equivalence classes of words M . Further, Schensted insertion is the inverse of w :

Given any tableau T , $w(T) = e$ for some word e . To see that $\emptyset \leftarrow e = T$, note that $w(\emptyset \leftarrow e) \equiv e$ by Proposition ???. That is, $e \equiv w(\emptyset \leftarrow e) \equiv w(T)$. Since a word is Knuth equivalent to the word of only one tableau, $\emptyset \leftarrow e = T$. On the other hand, if a word e gives rise to the tableau $T = \emptyset \leftarrow e$, then $w(T) = w(\emptyset \leftarrow e) \equiv e$.

This bijection between the associative monoid M and the set of tableaux implies that the set of tableaux forms a monoid under the multiplication defined on tableaux by the **insertion product**: $T \cdot U := T \leftarrow w(U)$. That is, for tableaux T, U, V :

- The empty tableau \emptyset is such that $T \cdot \emptyset = \emptyset \cdot T = T$.
- $T \cdot U$ is a tableau.
- $(T \cdot U) \cdot V = T \cdot (U \cdot V)$

6.1 Multiplication in the monoid

In general, the insertion product masks the outcome of multiplying two given tableaux, however there are circumstances under which some information is revealed. For example, given tableaux T and U , it is clear the the shape of $T \cdot U$ must contain the shape of T since each insertion adds boxes to the periphery of T . In the case that U is a tableau of row or column shape, even more can be said about the resulting shape of $T \cdot U$. Examine the following example:

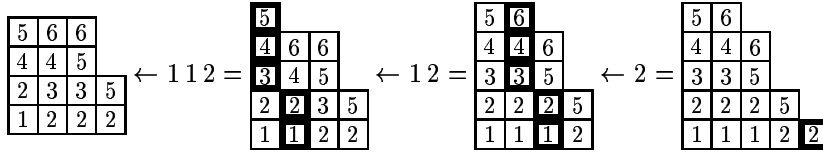
$$T \cdot U = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 4 & 5 \\ \hline 2 & 3 & 3 & 5 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array} \cdot \begin{array}{|c|c|c|} \hline 1 & 3 & 3 \\ \hline \end{array} = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 4 & 5 \\ \hline 2 & 3 & 3 & 5 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array} \leftarrow 133 = \begin{array}{|c|c|c|} \hline 5 & & \\ \hline 4 & 6 & 6 \\ \hline 3 & 4 & 5 \\ \hline 2 & 2 & 3 & 5 \\ \hline 1 & 1 & 2 & 2 \\ \hline \end{array} \leftarrow 33 = \begin{array}{|c|c|c|} \hline 5 & & \\ \hline 4 & 6 & 6 \\ \hline 3 & 4 & 5 \\ \hline 2 & 2 & 3 & 5 \\ \hline 1 & 1 & 2 & 2 & 3 \\ \hline \end{array} \leftarrow 3 = \begin{array}{|c|c|c|} \hline 5 & & \\ \hline 4 & 6 & 6 \\ \hline 3 & 4 & 5 \\ \hline 2 & 2 & 3 & 5 \\ \hline 1 & 1 & 2 & 2 & 3 & 3 \\ \hline \end{array}$$

Notice that in the process of multiplying $T \cdot U$, at most one cell has been added to any column of T . The following important lemma about Schensted insertion will lead to such properties on the product of tableaux.

²The proof is much more complex and is relegated to the Appendix

6.2 A result on bumping

The proof relies on the **bumping route**, the sequence of all cells affected during an insertion.

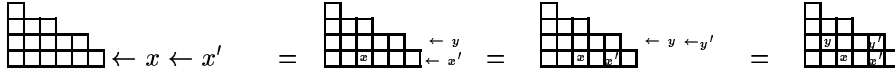


Observe that the bumping route of the first insertion lies to the left of that of the second.

The row bumping lemma: Insert x into the tableau T and denote the bumping route by R . Then insert x' into the resulting tableau and denote its bumping route by R' . Let c and c' denote the cells added to $T \leftarrow x$ and $T \leftarrow x \leftarrow x'$, respectively.

- If $x \leq x'$, then every cell of R is strictly left of those in R' and c is strictly left and weakly above c' .
- If $x > x'$, then every cell of R' is weakly left of those in R and c' is weakly left and strictly above c .

Proof. Consider the case that $x \leq x'$ bumps an element y from the first row. If an element y' is bumped by x' from the first row, then y' lies strictly to the right of y .



Repeating this argument gives that the bumping route R must lie strictly to the left of R' . The bumping route R ends at a cell c which must be strictly left and weakly above c' . Similarly we can show the second assertion. \square

Following from this lemma: If ν is the shape of $T \cdot U$ and λ is the shape of T . Let ν/λ denote the **skew diagram** obtained by removing the cells of λ from ν . Notice that no two cells lie in the same column. This holds in general, following from an important lemma about Schensted insertion.

Proposition 6. *If λ is the shape of T and ν is the shape of $T \cdot U$, then no two boxes in the skew diagram ν/λ lie in the same column if U is a row. Conversely, if X is a tableau of shape ν and $\lambda \subseteq \nu$ such that no two cells of ν/λ lie in the same column, then there is a tableau T of shape λ and a tableau U whose shape is a row such that $X = T \cdot U$.*

Proposition 7. *If the shape of U is a column, then no two boxes of ν/λ lie in the same row. Conversely, if X is a tableau of shape ν and $\lambda \subseteq \nu$ such that no two cells of ν/λ lie in the same row, then there is a tableau T of shape λ and a tableau U whose shape is a column such that $X = T \cdot U$.*

When U is a row, we say $\nu = \lambda +$ a horizontal n -strip where n is the number of cells in U and when U is a columns, $\nu = \lambda +$ a vertical n -strip.

6.3 The plactic algebra and Pieri rule

From any monoid, we can form a ring by considering the set of linear combinations of the elements in the monoid with coefficients in a fixed ring. Multiplication of two basis elements is determined by the monoid structure which then extends by bilinearity to all linear combinations. We will denote **the tableaux ring:** $R_{[m]}$ to be the ring associated to the tableau monoid.

- $R_{[m]}$ is a \mathbb{Z} -module with the tableaux in letters $\{1, \dots, m\}$ as a basis

- Multiplication in the ring can be defined by insertion product
- Associative, not commutative.

For ACE functions on words in the free monoid, see FREE.

- $\text{Free2PlaxClass}(w[.]);$ gives all the words in the same plactic class as $w[.]$ (i.e. all words knuth equivalent to $w[.]$).
- $\text{Free2Plax}(w[.]+w[...]+...);$ gives the decomposition of this sum into the plactic monoid. i.e. F/Knuth relations

In the tableau ring, we will be especially interested in $S_\lambda \in R_{[m]}$, the sum of all tableaux of shape λ . There are beautiful combinatorial formulas for the product $S_\lambda \cdot S_\mu$. For example, in the case that μ is a row or column we are equipped to prove

Theorem 8.

$$S_\lambda \cdot S_{(r)} = \sum_{\nu=\lambda+\text{horizontal } r\text{-strip}} S_\nu$$

$$S_\lambda \cdot S_{(1^r)} = \sum_{\nu=\lambda+\text{vertical } r\text{-strip}} S_\nu$$

Proof. The product of a tableau of shape λ and a row-tableau U gives a term occuring exactly once in one of the summands on the right. On the other hand, given any tableau X occuring in one of the S_ν , it can be written uniquely as $X = T \cdot U$ where T is a tableau of shape λ and U is a tableau whose shape is (r) . The second identity follows similarly. \square