

# Combinatorial Structures: Mombasa

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## 1 Permutations

A **permutation** of the integers  $\{1, \dots, n\}$  is a one-to-one map of the integers onto themselves:  $i \rightarrow \sigma_i$ . There are several ways to denote such an action.

**Two line notation:**  $\sigma = \begin{pmatrix} 1 & 2 & 3 & \cdots & n \\ \sigma_1 & \sigma_2 & \sigma_3 & \cdots & \sigma_n \end{pmatrix}$       **One line notation:**  $\sigma = (\sigma_1 \sigma_2 \dots \sigma_n)$

The set of all permutations of the integers  $\{1, \dots, n\}$  forms a group called the **symmetric group**,  $S_n$ , under composition of maps.

$$S_3 = \{(123), (132), (312), (213), (231), (321)\}$$

$$\tau = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix} \text{ and } \sigma = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} \implies \tau\sigma = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 3 & 2 \end{pmatrix}$$

The symmetric group is not an abelian group since the elements do not necessarily commute.

$$\sigma\tau = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix}$$

The **inverse** of  $\sigma \in S_n$  is the element  $\sigma^{-1}$  where  $\sigma^{-1}\sigma = id$ . Since the inverse of a permutation  $\sigma$  is the permutation that undoes the action of  $\sigma$ , it can easily be computed by:

$$\sigma = \begin{pmatrix} 1 & 2 & 3 & \cdots & n \\ \sigma_1 & \sigma_2 & \sigma_3 & \cdots & \sigma_n \end{pmatrix} \implies \sigma^{-1} = \begin{pmatrix} \sigma_1 & \sigma_2 & \sigma_3 & \cdots & \sigma_n \\ 1 & 2 & 3 & \cdots & n \end{pmatrix} = \begin{pmatrix} 1 & 2 & 3 & \cdots & n \\ \sigma'_1 & \sigma'_2 & \sigma'_3 & \cdots & \sigma'_n \end{pmatrix}$$

$$\sigma = (32514) \implies \sigma^{-1} = \begin{pmatrix} 3 & 2 & 5 & 1 & 4 \\ 1 & 2 & 3 & 4 & 5 \end{pmatrix} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 4 & 2 & 1 & 5 & 3 \end{pmatrix}$$

A permutation that is equal to its inverse is called an **involution**. Note  $\sigma = \sigma^{-1} \iff \sigma\sigma = id$ . The involutions of  $S_3$  are  $(123)$ ,  $(213)$ ,  $(321)$ ,  $(132)$ . For example, check  $(213)$ :

$$\begin{pmatrix} 2 & 1 & 3 \\ 1 & 2 & 3 \end{pmatrix} \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \end{pmatrix}$$

There is a third notation used in the study of permutations called **cycle notation**. In this case,  $\sigma = (a_1, \dots, a_\ell) \cdots (b_1, \dots, b_j)$  means  $\sigma(a_i) = a_{i+1}$  and  $\sigma(a_\ell) = a_1, \dots, \sigma(b_i) = b_{i+1}$  and  $\sigma(b_j) = b_1$ .

$$\sigma = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} = (213) = (1,2)(3)$$

The **cycle structure** of a permutation  $\sigma \in S_n$  is the non-increasing vector of integers  $\lambda = (\lambda_1, \dots, \lambda_\ell)$  determined by the number of elements in the disjoint cycles of  $\sigma$ . For example,

$$\text{one-line notation} = \{(123), (213), (312), (231), (321), (132)\}$$

$$\text{cycle notation} = \{(1)(2)(3), (1,2)(3), (1,3,2), (1,2,3), (1,3)(2), (1)(2,3)\}$$

$$\text{cycle structures} = \{(1,1,1), (2,1), (3), (3), (2,1), (2,1)\}$$

Recall that the conjugacy class of a permutation  $\sigma$  is the set of all permutations,  $\{\gamma^{-1}\sigma\gamma : \gamma \in S_n\}$ . It turns out that the conjugacy class of a given permutation can be determined by finding all permutation with the same cycle structure. For instance, since the cycle structure of  $(321)$  is  $(2,1)$ , the elements in the conjugacy class of  $(321)$  are  $(321)$ ,  $(132)$ , and  $(213)$ .

## 2 Partitions

The conjugacy classes are one of many examples of objects that can be indexed by weakly decreasing vectors of positive integers. This is one of many reasons to closely study such vectors. These vectors are called partitions. That is, a **partition** is a vector of integers  $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_n)$  where  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n > 0$ . The partitions of 4 are

$$(4), (3,1), (2,2), (2,1,1), (1,1,1,1)$$

We can express partitions in an alternative notation that emphasizes the number of times a given part occurs. That is,  $\lambda = (\lambda_1^{m_1}, \lambda_2^{m_2}, \dots, \lambda_\ell^{m_\ell})$  where there are  $m_i$  occurrences of  $\lambda_i$ . In this case, the partitions of 4 are denoted

$$(4), (3,1), (2^2), (2,1^2), (1^4)$$





## 4 Robinson-Schensted Correspondence

### 4.1 Insertion/deletion algorithm

A fundamental operation on tableaux is given by the **Schensted insertion algorithm**, which sends a tableau  $T$  and a positive integer  $x$  to a new tableau that has one more box than  $T$ . This new tableau is denoted  $T \leftarrow x$ .

If  $x$  is weakly larger than every entry in the bottom row of  $T$ , add  $x$  to the end of this row. If not, find the leftmost entry  $e$  in the bottom row that is strictly larger than  $x$ . Replace this entry by  $x$ . Repeat the process on the next row with the letter  $e$ .

$$\begin{array}{|c|c|c|} \hline 5 & & \\ \hline 3 & 3 & 4 \\ \hline 1 & 1 & 2 & 3 \\ \hline \end{array} \leftarrow 3 = \begin{array}{|c|c|c|} \hline 5 & & \\ \hline 3 & 3 & 4 \\ \hline 1 & 1 & 2 & 3 & 3 \\ \hline \end{array}$$

$$\begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 5 & \\ \hline 1 & 1 & 3 & 4 \\ \hline \end{array} \leftarrow 3 = \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 5 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \leftarrow 4 = \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \leftarrow 5 = \begin{array}{|c|c|c|} \hline 3 & 5 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \leftarrow 6 = \begin{array}{|c|c|c|} \hline 6 & & \\ \hline 3 & 5 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array}$$

Note that this algorithm will send any word to a tableau by successively inserting the letters using Schensted's insertion algorithm.

$$\emptyset \leftarrow 234412232 = \begin{array}{|c|} \hline 2 \\ \hline \end{array} \leftarrow 34412232 = \begin{array}{|c|c|} \hline 2 & 3 \\ \hline \end{array} \leftarrow 4412232 = \dots = \begin{array}{|c|c|c|c|} \hline 2 & 3 & 4 & 4 \\ \hline \end{array} \leftarrow 12232$$

$$= \begin{array}{|c|c|c|c|} \hline 2 & 3 & 4 & 4 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array} \leftarrow 12232 = \begin{array}{|c|c|c|} \hline 2 & 3 & \\ \hline 1 & 2 & 4 & 4 \\ \hline \end{array} \leftarrow 232 = \begin{array}{|c|c|c|} \hline 2 & 3 & 4 \\ \hline 1 & 2 & 2 & 4 \\ \hline \end{array} \leftarrow 32 = \begin{array}{|c|c|c|c|} \hline 2 & 3 & 4 & 4 \\ \hline 1 & 2 & 2 & 3 \\ \hline \end{array} \leftarrow 2 = \begin{array}{|c|c|c|c|} \hline 4 & & & \\ \hline 2 & 3 & 3 & 4 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array}$$

Given that  $\hat{T} = T \leftarrow x$  has a new cell in position  $c$ , we can easily reverse the process of insertion starting from the entry in position  $c$ . The **deletion algorithm** will send a given tableau  $T$  and a corner  $c$  of  $T$  to a tableau of the shape of  $T$  with the cell  $c$  removed.

Delete the entry  $x$  from corner  $c$ . If corner  $c$  was in the bottom row of  $T$ , we are done. If not, find the rightmost entry  $e$  smaller than  $x$  in the row below that with  $c$  and replace  $e$  by  $x$ . Repeat this process with entry  $e$ .

If  $c$  is the cell containing 6 in the tableau  $T$ :

$$T = \begin{array}{|c|c|c|} \hline 6 & & \\ \hline 3 & 5 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|} \hline 3 & 5 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \leftarrow 6 \rightarrow \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 4 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \leftarrow 5 \rightarrow \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 5 & \\ \hline 1 & 1 & 3 & 3 \\ \hline \end{array} \leftarrow 4 \rightarrow \begin{array}{|c|c|c|} \hline 3 & 6 & \\ \hline 2 & 5 & \\ \hline 1 & 1 & 3 & 4 \\ \hline \end{array}$$

### 4.2 The correspondence

A bijection between permutations and pairs of standard tableaux with the same shape arises from the insertion algorithm. The **Robinson-Schensted algorithm** is a method for constructing a pair of tableaux of the same shape from a permutation.

From any  $\sigma \in S_n$ , the first tableau in our pair is constructed from  $\sigma$  using the Schensted insertion algorithm. However, during this construction when a new box is added by inserting  $\sigma_i$ , put the letter  $i$  into this box in the second tableau. The first tableau is called the **insertion tableau** and the second is the **recording tableau**.

To find the pair of tableaux associated to  $\sigma = 415326$ , start by row inserting  $\sigma_1 = 4$  into  $\emptyset$ ;

$$\emptyset \leftarrow 4 = \begin{array}{|c|} \hline 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 1 \\ \hline \end{array}$$

$$\begin{array}{|c|} \hline 4 \\ \hline \end{array} \leftarrow 1 = \begin{array}{|c|} \hline 4 \\ \hline 1 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 2 \\ \hline 1 \\ \hline \end{array}$$

$$\begin{array}{l}
\begin{array}{|c|} \hline 4 \\ \hline 1 \\ \hline \end{array} \leftarrow 5 = \begin{array}{|c|c|} \hline 4 & \\ \hline 1 & 5 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 2 \\ \hline 1 \\ \hline 3 \\ \hline \end{array} \\
\begin{array}{|c|c|} \hline 4 & \\ \hline 1 & 5 \\ \hline \end{array} \leftarrow 3 = \begin{array}{|c|c|} \hline 4 & 5 \\ \hline 1 & 3 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 2 & 4 \\ \hline 1 & 3 \\ \hline \end{array} \\
\begin{array}{|c|c|} \hline 4 & 5 \\ \hline 1 & 3 \\ \hline \end{array} \leftarrow 2 = \begin{array}{|c|c|} \hline 4 & \\ \hline 3 & 5 \\ \hline 1 & 2 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 5 & \\ \hline 2 & 4 \\ \hline 1 & 3 \\ \hline \end{array} \\
\begin{array}{|c|c|} \hline 4 & \\ \hline 3 & 5 \\ \hline 1 & 2 \\ \hline \end{array} \leftarrow 6 = \begin{array}{|c|c|} \hline 4 & \\ \hline 3 & 5 \\ \hline 1 & 2 & 6 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 5 & \\ \hline 2 & 4 \\ \hline 1 & 3 & 6 \\ \hline \end{array}
\end{array}$$

Note that the recording tableau is a standard tableau of  $n$  since  $Q$  is formed by adding elements in increasing order to cells on the periphery.

Since the Schensted insertion algorithm is reversible, this process can be inverted to obtain a permutation of  $S_n$  from a pair of standard tableaux of the same shape. That is, starting with the entry  $e$  in the insertion tableau that lies in position  $n$  determined by the recording tableaux, we use the deletion process to remove this letter. Then repeat starting with  $n - 1$ . Thus, the **Robinson-Schensted Correspondence** is a bijection between elements of  $S_n$  and the set of ordered pairs  $(P, Q)$  of standard tableaux of  $n$  having the same shape.

One direct consequence of the RS-correspondence is that:

**Theorem 1.** *The number of permutations of  $\{1, 2, \dots, n\}$  is equal to the number of pairs of standard tableaux of the same shape  $\lambda$  as  $\lambda$  varies over all partitions of  $n$ .*

$$\sum_{\lambda \vdash n} (f^\lambda)^2 = n!$$

When we consider only permutation that are involutions, another beautiful identity holds. Notice the number of involutions of  $S_4$ :

$$(1\ 2\ 3\ 4), (2\ 1\ 3\ 4), (3\ 2\ 1\ 4), (4\ 2\ 3\ 1), (1\ 3\ 2\ 4), (1\ 4\ 3\ 2), (1\ 2\ 4\ 3), (2\ 1\ 4\ 3), (3\ 4\ 1\ 2), (4\ 3\ 2\ 1)$$

equals the number of standard tableaux of degree 4:

$$\begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 2 \\ \hline 1 & 3 & 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 3 \\ \hline 2 \\ \hline 1 & 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 4 \\ \hline 2 \\ \hline 1 & 3 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 3 \\ \hline 1 & 2 & 4 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 4 \\ \hline 3 \\ \hline 1 & 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 4 \\ \hline 1 & 2 & 3 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 2 & 4 \\ \hline 1 & 3 \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline 3 & 4 \\ \hline 1 & 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 4 \\ \hline 3 \\ \hline 2 \\ \hline 1 \\ \hline \end{array}.$$

This turns out to be true in general:

**Theorem 2.** *The number of involutions on  $\{1, \dots, n\}$  is the number of standard tableaux on  $n$  letters.*

The argument relies on an interesting result: If  $\sigma$  corresponds to the pair of standard tableaux  $(P, Q)$  under the RS-correspondence, then  $\sigma^{-1}$  corresponds to the pair  $(Q, P)$ . Thus for  $\sigma = \sigma^{-1}$  an involution,  $(P, Q) = (Q, P)$ . That is,  $\sigma$  corresponds to a pair of the same standard tableaux.

## 5 Plactic monoid

Another important idea linked closely to Schensted insertion is the study of tableaux as a monoid. Recall that an associative **monoid** is a set  $G$  containing an identity  $e$  ( $x \cdot e = e \cdot x = x$  for all  $x \in G$ ) and having a binary operation (called multiplication) for which the associative law holds. That is, a monoid is a group without an inverse.

The collection of all words on a given alphabet forms a monoid called the **free monoid**, where multiplication is defined by word juxtaposition:

$$w = (3416647) \quad \text{and} \quad w' = (2155) \quad \text{give} \quad ww' = (34166472155)$$

The empty word  $\emptyset$  is the identity element. In particular, we have the monoid  $F$  on letters  $\{1, \dots, m\}$ . Using the obvious structure of the free monoid, a connection between words and tableaux that will enable us to induce a monoid on tableaux.

## 5.1 Knuth equivalence

The **row word** of tableau  $T$ , denoted  $w(T)$ , is the sequence of letters read from a tableau like a book in English.

$$w\left(\begin{array}{|c|c|} \hline 3 & 5 \\ \hline 2 & 4 \\ \hline 1 & 1 & 2 & 2 & 6 \\ \hline \end{array}\right) = 35241126$$

Note that **every tableau**  $T \rightarrow$  **a word**  $w(T)$ . The original tableau is recovered from its row word by noting that the rows end at each **descent** – any position in a word  $w$  where  $w_i > w_{i+1}$ :

$$35|24|1126 \longrightarrow \begin{array}{|c|c|} \hline 3 & 5 \\ \hline 2 & 4 \\ \hline 1 & 1 & 2 & 2 & 6 \\ \hline \end{array}$$

However, not every word comes from a tableaux. Only when a word that is cut at its descents has pieces of weakly increasing length, and when stacked have columns that are strictly increasing does the word form a tableau. Otherwise, we may find the following scenario:

$$357241126 \rightarrow 357|24|1126 \rightarrow \begin{array}{|c|c|c|} \hline 3 & 5 & 7 \\ \hline 2 & 4 & \\ \hline 1 & 1 & 2 & 2 & 6 \\ \hline \end{array}$$

*Some words*  $\rightarrow$  **a tableau**.

To reconcile that there are “too many” words, we shall classify certain words as “the same”. The **Knuth relation**  $A$  is move on a three letter word:

$$y z x \Rightarrow y x z \iff \boxed{y} \boxed{z} \leftarrow x = \begin{array}{|c|} \hline y \\ \hline x \\ \hline \end{array} \boxed{z} \quad \text{if } x < y \leq z$$

The **Knuth relation**  $B$  is move on a three letter word:

$$x z y \Rightarrow z x y \iff \boxed{x} \boxed{z} \leftarrow y = \begin{array}{|c|} \hline z \\ \hline x \\ \hline \end{array} \boxed{y} \quad \text{if } x \leq y < z$$

For example,

$$342 \stackrel{A}{\equiv} 324$$

An **elementary Knuth transformation** on a word applies one of  $A$  or  $B$  or their inverses to three successive letters of a word. Thus, words  $w$  and  $w'$  are **Knuth equivalent** if  $w$  can be obtained by applying a sequence of elementary Knuth transformations to  $w'$ . For example,

$$1223324 \equiv 1232234 \quad \text{since} \quad 1223324 \stackrel{A}{\equiv} 1223234 \stackrel{B}{\equiv} 1232234$$

The set  $M = F/R$  where  $R$  is the equivalence relation generated by the Knuth relations (the Knuth equivalence classes of words) is closed under multiplication (juxtaposition). That is,  $w, u \in M$  are such that  $wu \equiv w'u'$  if  $w \equiv w'$  and  $u \equiv u'$ . Thus the multiplication on  $F$  descends to a multiplication on the set  $M$ . This makes  $M$  into an associative monoid called the **plactic monoid**.

It will develop that **each equivalence class**  $\rightarrow$  **tableau**. Then, a bijection between equivalence classes and tableaux will reveal that tableaux form a monoid that is isomorphic to the plactic monoid.

## 5.2 Knuth slow insertion

Recall, the insertion algorithm produces a tableau  $T \leftarrow x$  from a tableau  $T$  and letter  $x$ . This algorithm can instead be achieved systematically by applying a sequence of Knuth transformations  $A$  and  $B$  to the left of  $w(T)x$ .

Consider only the row word of the bottom row of  $T$ ,  $w = u_1 \cdots u_p x' v_1 \cdots v_q x$  where  $x' > x \geq u_p$ . Check that  $x$  is strictly smaller than its two preceding letters. If so, we move  $x$  forward by transposing with the entry to its right since  $v_q \geq v_{q-1} > x$  implies the Knuth operation  $A$  moves  $x$  to the left of  $v_q$ .

$$122334\underline{2} \stackrel{A}{\equiv} 12233\underline{2}4 \stackrel{A}{\equiv} 1223\underline{2}34$$

This continues until  $x$  is no longer smaller, then  $x$  rests while  $x'$  is pushed forward by transposing with the entry to its left. That is, use transformation  $A$  until the configuration  $u_1 \cdots u_p x' x v_1 \cdots v_q$  is reached. Then  $u_p \leq x$  and  $x' > x$  implies that operation  $B$  moves  $x'$  to the left of  $u_{p-1}$ .

$$1223\underline{2}34 \stackrel{B}{\equiv} 123\underline{2}234 \stackrel{B}{\equiv} 13\underline{2}2234 \stackrel{B}{\equiv} \underline{3}122234$$

This continues until  $x'$  has been moved to the beginning of the row word  $w$  of the first row of  $T$ . As such, the letter  $x$  has been Schensted inserted into the bottom row and  $x'$  has been bumped. This process is then repeated with  $x'$  at the end of the row word of the second row of  $T$ , amounting exactly to the insertion of  $x'$  into the second row, etc. Thus, the Schensted insertion of  $x$  into  $T$  can be achieved by Knuth transformations on  $w(T)x$ . That is,

**Proposition 3.** *For any tableau  $T$  and positive integer  $x$ ,*

$$w(T \leftarrow x) \equiv w(T)x$$

More generally, given any word  $w = x_1 \cdots x_\ell$  and starting from the empty tableau  $T = \emptyset$  we have  $w(\emptyset \leftarrow x_1) \equiv x_1$ , then with  $T = \emptyset \leftarrow x_1$ ,  $w(T \leftarrow x_2) \equiv x_1 x_2$ , and so. Thus proving:

**Proposition 4.** *Every word is Knuth equivalent to the word of a tableau. In particular,*

$$x_1 x_2 \cdots x_\ell \equiv w(\emptyset \leftarrow x_1 \leftarrow \cdots \leftarrow x_\ell)$$

implying that the word  $w = x_1 \cdots x_\ell$  is Knuth equivalent to the tableau  $T = \emptyset \leftarrow w$ .

In fact a stronger result holds<sup>1</sup>. ...

**Theorem 5.** *Every word is Knuth equivalent to the word of a unique tableau. That is, each equivalence class contains exactly one row word of a tableau.*

All words that are Knuth equivalent can be associated to the same tableau and no others:

$$\begin{aligned} 235512232 &\equiv 235512322 \equiv 235513222 \equiv 235531222 \\ &\equiv 235351222 \equiv 253351222 \equiv 523351222 \equiv w\left(\begin{array}{|c|c|c|c|} \hline 5 & & & \\ \hline 2 & 3 & 3 & 5 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array}\right) \end{aligned}$$

## 6 The tableaux monoid and the plactic algebra

We have now seen that the map  $w$  from a tableau to the row word of this tableau is a bijection between the set of tableaux and the Knuth equivalence classes of words  $M$ . Further, Schensted insertion is the inverse of  $w$ :

Given any tableau  $T$ ,  $w(T) = e$  for some word  $e$ . To see that  $\emptyset \leftarrow e = T$ , note that  $w(\emptyset \leftarrow e) \equiv e$  by Proposition ???. That is,  $e \equiv w(\emptyset \leftarrow e) \equiv w(T)$ . Since a word is Knuth equivalent to the word of only one tableaux,  $\emptyset \leftarrow e = T$ . On the other hand, if a word  $e$  gives rise to the tableau  $T = \emptyset \leftarrow e$ , then  $w(T) = w(\emptyset \leftarrow e) \equiv e$ .

<sup>1</sup>The proof is much more complex and is relegated to the Appendix

This bijection between the associative monoid  $M$  and the set of tableaux implies that the set of tableaux forms a monoid under the multiplication defined on tableaux by the **insertion product**:  $T \cdot U := T \leftarrow w(U)$ . That is, for tableaux  $T, U, V$ :

- The empty tableau  $\emptyset$  is such that  $T \cdot \emptyset = \emptyset \cdot T = T$ .
- $T \cdot U$  is a tableau.
- $(T \cdot U) \cdot V = T \cdot (U \cdot V)$

## 6.1 Multiplication in the monoid

In general, the insertion product masks the outcome of multiplying two given tableaux, however there are circumstances under which some information is revealed. For example, given tableaux  $T$  and  $U$ , it is clear the the shape of  $T \cdot U$  must contain the shape of  $T$  since each insertion adds boxes to the periphery of  $T$ . In the case that  $U$  is a tableau of row or column shape, even more can be said about the resulting shape of  $T \cdot U$ . Examine the following example:

$$T \cdot U = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 4 & 5 \\ \hline 2 & 3 & 3 & 5 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array} \cdot \begin{array}{|c|c|c|} \hline 1 & 1 & 3 & 3 \\ \hline \end{array} = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 4 & 5 \\ \hline 2 & 3 & 3 & 5 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array} \leftarrow 1333 = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 6 & 6 \\ \hline 3 & 4 & 5 \\ \hline 2 & 2 & 3 & 5 \\ \hline 1 & 1 & 2 & 2 \\ \hline \end{array} \leftarrow 33 = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 6 & 6 \\ \hline 3 & 4 & 5 \\ \hline 2 & 2 & 3 & 5 \\ \hline 1 & 1 & 2 & 2 & 3 \\ \hline \end{array} \leftarrow 3 = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 6 & 6 \\ \hline 3 & 4 & 5 \\ \hline 2 & 2 & 3 & 5 \\ \hline 1 & 1 & 2 & 2 & 3 & 3 \\ \hline \end{array}$$

Notice that in the process of multiplying  $T \cdot U$ , at most one cell has been added to any column of  $T$ . The following important lemma about Schensted insertion will lead to such properties on the product of tableaux.

## 6.2 A result on bumping

The proof relies on the **bumping route**, the sequence of all cells affected during an insertion.

$$\begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 4 & 5 \\ \hline 2 & 3 & 3 & 5 \\ \hline 1 & 2 & 2 & 2 \\ \hline \end{array} \leftarrow 112 = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 6 & 6 \\ \hline 3 & 4 & 5 \\ \hline 2 & 2 & 3 & 5 \\ \hline 1 & 1 & 2 & 2 \\ \hline \end{array} \leftarrow 12 = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 4 & 6 \\ \hline 3 & 3 & 5 \\ \hline 2 & 2 & 2 & 5 \\ \hline 1 & 1 & 1 & 2 \\ \hline \end{array} \leftarrow 2 = \begin{array}{|c|c|c|} \hline 5 & 6 & 6 \\ \hline 4 & 4 & 6 \\ \hline 3 & 3 & 5 \\ \hline 2 & 2 & 2 & 5 \\ \hline 1 & 1 & 1 & 2 & 2 \\ \hline \end{array}$$

Observe that the bumping route of the first insertion lies to the left of that of the second.

**The row bumping lemma:** Insert  $x$  into the tableau  $T$  and denote the bumping route by  $R$ . Then insert  $x'$  into the resulting tableau and denote its bumping route by  $R'$ . Let  $c$  and  $c'$  denote the cells added to  $T \leftarrow x$  and  $T \leftarrow x \leftarrow x'$ , respectively.

- If  $x \leq x'$ , then every cell of  $R$  is strictly left of those in  $R'$  and  $c$  is strictly left and weakly above  $c'$ .
- If  $x > x'$ , then every cell of  $R'$  is weakly left of those in  $R$  and  $c'$  is weakly left and strictly above  $c$ .

*Proof.* Consider the case that  $x \leq x'$  bumps an element  $y$  from the first row. If an element  $y'$  is bumped by  $x'$  from the first row, then  $y'$  lies strictly to the right of  $y$ .

$$\begin{array}{|c|c|c|c|} \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \leftarrow x \leftarrow x' = \begin{array}{|c|c|c|c|} \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \leftarrow y \leftarrow x' = \begin{array}{|c|c|c|c|} \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \leftarrow y \leftarrow y' = \begin{array}{|c|c|c|c|} \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline \end{array}$$

Repeating this argument gives that the bumping route  $R$  must lie strictly to the left of  $R'$ . The bumping route  $R$  ends at a cell  $c$  which must be strictly left and weakly above  $c'$ . Similarly we can show the second assertion.  $\square$

Following from this lemma: If  $\nu$  is the shape of  $T \cdot U$  and  $\lambda$  is the shape of  $T$ . Let  $\nu/\lambda$  denote the **skew diagram** obtained by removing the cells of  $\lambda$  from  $\nu$ . Notice that no two cells lie in the same column. This holds in general, following from an important lemma about Schensted insertion.

**Proposition 6.** *If  $\lambda$  is the shape of  $T$  and  $\nu$  is the shape of  $T \cdot U$ , then no two boxes in the skew diagram  $\nu/\lambda$  lie in the same column if  $U$  is a row. Conversely, if  $X$  is a tableau of shape  $\nu$  and  $\lambda \subseteq \nu$  such that no two cells of  $\nu/\lambda$  lie in the same column, then there is a tableau  $T$  of shape  $\lambda$  and a tableau  $U$  whose shape is a row such that  $X = T \cdot U$ .*

**Proposition 7.** *If the shape of  $U$  is a column, then no two boxes of  $\nu/\lambda$  lie in the same row. Conversely, if  $X$  is a tableau of shape  $\nu$  and  $\lambda \subseteq \nu$  such that no two cells of  $\nu/\lambda$  lie in the same row, then there is a tableau  $T$  of shape  $\lambda$  and a tableau  $U$  whose shape is a column such that  $X = T \cdot U$ .*

When  $U$  is a row, we say  $\nu = \lambda +$  a horizontal  $n$ -strip where  $n$  is the number of cells in  $U$  and when  $U$  is a columns,  $\nu = \lambda +$  a vertical  $n$ -strip.

### 6.3 The plactic algebra and Pieri rule

From any monoid, we can form a ring by considering the set of linear combinations of the elements in the monoid with coefficients in a fixed ring. Multiplication of two basis elements is determined by the monoid structure which then extends by bilinearity to all linear combinations. We will denote **the tableaux ring**:  $R_{[m]}$  to be the ring associated to the tableau monoid.

- $R_{[m]}$  is a  $\mathbb{Z}$ -module with the tableaux in letters  $\{1, \dots, m\}$  as a basis
- Multiplication in the ring can be defined by insertion product
- Associative, not commutative.

In the tableau ring, we will be especially interested in  $S_\lambda \in R_{[m]}$ , the sum of all tableaux of shape  $\lambda$ . There are beautiful combinatorial formulas for the product  $S_\lambda \cdot S_\mu$ . For example, in the case that  $\mu$  is a row or column we are equipped to prove

**Theorem 8.**

$$S_\lambda \cdot S_{(r)} = \sum_{\nu=\lambda+\text{horizontal } r\text{-strip}} S_\nu$$

$$S_\lambda \cdot S_{(1^r)} = \sum_{\nu=\lambda+\text{vertical } r\text{-strip}} S_\nu$$

*Proof.* The product of a tableau of shape  $\lambda$  and a row-tableau  $U$  gives a term occurring exactly once in one of the summands on the right. On the other hand, given any tableau  $X$  occurring in one of the  $S_\nu$ , it can be written uniquely as  $X = T \cdot U$  where  $T$  is a tableau of shape  $\lambda$  and  $U$  is a tableau whose shape is  $(r)$ . The second identity follows similarly.  $\square$

## 7 Appendix

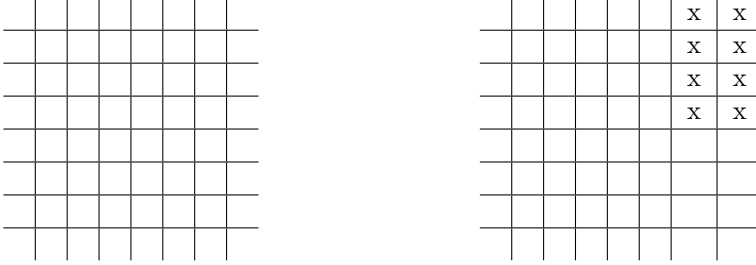
### 7.1 Viennot's construction

Viennot's construction is an alternative method for obtaining the RS-insertion and recording tableaux from a given permutation. This method has the advantage of clearly showing that if  $\sigma$  corresponds to the pair of standard tableaux  $(P, Q)$  under the RS-correspondence, then  $\sigma^{-1}$  corresponds to the pair  $(Q, P)$ .

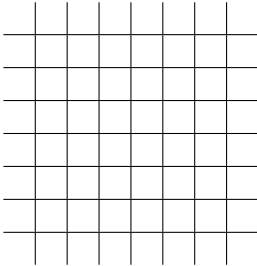
We start by identifying each permutation  $\sigma = (\sigma_1 \cdots \sigma_n)$  with the set of lattice points in  $\mathbb{N} \times \mathbb{N}$ :

$$\{(i, \sigma_i) : i = 1, \dots, n\}$$

The *shadow* of any lattice point  $(a, b)$  is the region to the right of  $x = a$  and above  $y = b$ . For example with  $\sigma = (6\ 2\ 3\ 1\ 5\ 4\ 7)$ , we obtain:



The shaded region depicts the shadow of  $(6, 4)$ . The *special points* in a set of lattice points  $\mathcal{S}$  are those that do not lie in the shadow of any other point in  $\mathcal{S}$ . For example, the special points in the set  $\mathcal{S} = \{(i, \sigma) : i = 1, \dots, n\}$  are  $\{(1, 6), (2, 2), (4, 1)\}$ . Given a set of lattice points  $\mathcal{S}$ , the *shadow line* of  $\mathcal{S}$  is the boundary of the region obtained by casting shadows of the special points. Taking all the points in set  $\mathcal{S}$ , the shadow line is:



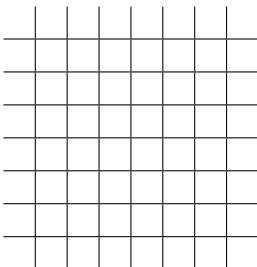
Since every shadow is cast upward and to the right, a shadow line must be a series of vertical and horizontal steps moving to the right and downward.

Furthermore,

*Remark 9.* Lower corners  $\lfloor$  in the shadow line of  $\mathcal{S}$  are determined by the special points of  $\mathcal{S}$ .

*Remark 10.* If the shadow line is horizontal at  $y = \sigma_i$  then it remains horizontal until reaching an  $x$ -coordinate where  $\sigma_x < \sigma_i$ .

We shall consider the set of shadow lines  $\{L_1, L_2, \dots\}$ , where  $L_1$  is the shadow line obtained from the set  $\mathcal{S}$  of all points determined by  $\sigma$ .  $L_2$  is the shadow line of  $\mathcal{S}^2$  – the set of points in  $\mathcal{S}$  minus the special points. In this manner,  $L_i$  is the shadow line of the set  $\mathcal{S}^i$  obtained by deleting the special points from  $\mathcal{S}^{i-1}$ . The *shadow diagram* of  $\sigma$  is the set of shadow lines  $\{L_1, L_2, \dots\}$ . In our example, the shadow diagram is:





In the shadow diagram of  $\sigma$ ,  $\sigma_{i+1} > y_\ell > \cdots > y_1$  implies that the lowest points on all shadow lines are lower than  $y = \sigma_{i+1}$  and thus a new shadow line begins at  $x = i + 1$  with bottom corner  $(i + 1, \sigma_{i+1})$ .

$$\begin{array}{r}
 \cdot - \sigma_{i+1} \\
 - - y_\ell \\
 - - \cdot \\
 | - \vdots \\
 - - \vdots \\
 - - y_2 \\
 - - y_1
 \end{array}$$

Therefore, the bottom row of  $P_{i+1}$  contains the lowest  $y$ -coordinates of shadow lines at  $x = i + 1$ .

On the other hand, if there is some  $j$  where  $y_1 < \cdots < y_{j-1} < \sigma_{i+1} < y_j$ , then row insertion implies that the entry  $y_j$  is replaced by  $\sigma_{i+1}$  in the bottom row of  $P_{i+1} = P_i \leftarrow \sigma_{i+1}$ .

$$P^{i+1} = \begin{array}{|c|c|c|c|c|c|c|c|}
 \hline
 & & & & & & & \\
 \hline
 & & & & \leftarrow y_j & & & \\
 \hline
 y_1 & \cdots & y_{j-1} & \sigma_{i+1} & y_{j+1} & & \cdots & y_\ell \\
 \hline
 \end{array}$$

Note that since a cell has not been added to the bottom row,  $i + 1$  does not occur in the bottom row of  $Q$ .

In the shadow diagram when  $y_{j-1} < \sigma_{i+1} < y_j$ ,  $(i + 1, \sigma_{i+1})$  lies on the shadow line that was at height  $y_j$  when  $x = i$ . Thus, the lowest coordinate  $y_j$  of this line is replaced by  $\sigma_{i+1}$  (depicted by a vertical move) while all others remain at the same height (since they are horizontal moves).

$$\begin{array}{r}
 - - y_\ell \\
 - - \vdots \\
 - - y_j \\
 \quad - \sigma_{i+1} \\
 - - y_{j-1} \\
 - - \vdots \\
 - - y_1
 \end{array}$$

Therefore, the bottom row of  $P_{i+1}$  contains exactly the lowest  $y$ -coordinates of the shadow lines at  $x = i + 1$  as claimed.

In summary, the shadow diagram of  $\sigma$  depicts the changes to the first row of the insertion tableau of  $\emptyset \leftarrow \sigma$ . That is, a vertical segment at  $x = i$  indicates that a letter is bumped from the bottom row when  $\sigma_i$  is inserted, and if the vertical piece starts a new shadow line, nothing is bumped and  $\sigma_i$  is added to the end of the first row.

Since the recording tableau puts letter  $i$  in the cell added to  $P_{i-1}$  after inserting  $\sigma_i$ , and a cell is added to the bottom row only when a new shadow line is started, the bottom row of the recording tableau  $Q$  is determined by the coordinates  $x_1, \dots, x_n$  that mark the start of the shadow lines  $L_1, \dots, L_n$ . □

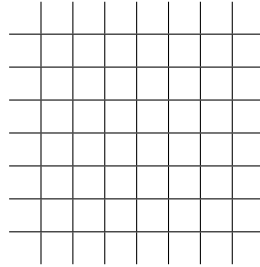
In particular, if we examine the vertical line  $x = n$ , we find that the bottom row of  $P = \emptyset \leftarrow \sigma$  is given by the lowest coordinates  $y_1, \dots, y_n$  of the shadow lines  $L_1, \dots, L_n$ . And if we examine the horizontal line  $y = n$ , we find that the bottom row of  $Q$  is given by the leftmost  $x$ -coordinates of the shadow lines  $L_1, \dots, L_n$ .

If  $i \rightarrow \sigma_i$  under  $\sigma$  then  $\sigma_i \rightarrow i$  under  $\sigma^{-1}$ . Thus, the shadow diagram for  $\sigma^{-1}$  is simply the reflection of the shadow diagram for  $\sigma$  about the line  $y = x$  (i.e.  $(i, \sigma_i) \rightarrow (\sigma_i, i)$ ). Since we can

obtain the bottom row of  $P$  and  $Q$  in  $\sigma \leftrightarrow (P, Q)$  by reading the appropriate  $x$  and  $y$ -coordinates from the shadow diagram for  $\sigma$ , it is thus clear that with  $\sigma^{-1} \leftrightarrow (\bar{P}, \bar{Q})$ , the bottom row of  $\bar{P}$  is given by the bottom row of  $Q$  and the bottom row of  $\bar{Q}$  is given by the bottom row of  $P$ .

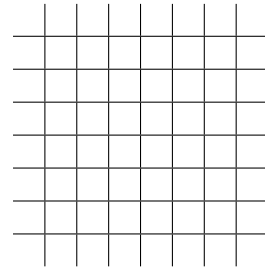
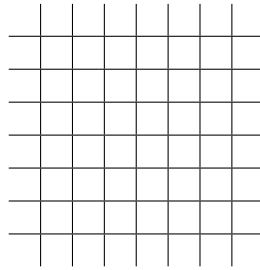
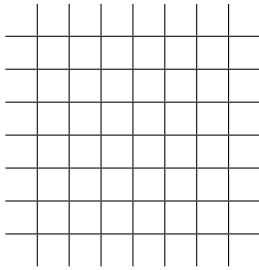
To construct the remainder of the tableaux  $(P, Q)$  from the shadow diagram of  $\sigma$ , note that if an upper corner on any shadow diagram is the point  $(i, a)$ , then the letter  $a$  was bumped from the

first row of  $P^i$  with the insertion of  $\sigma_{i+1}$ .



Thus, if  $y_1, \dots, y_r$  denote the set of

$y$ -coordinates of the upper corners in the shadow diagram of  $\sigma$  (where  $y_1$  is the leftmost such corner and  $y_r$  the rightmost), the tableau  $P$  (minus the bottom row) is exactly  $\emptyset \leftarrow y_1 \leftarrow \dots \leftarrow y_r$ .



If we construct a new shadow diagram using the set of upper corners, we can apply the previous result to obtain the second row of  $P$  and  $Q$ . By iterating this procedure, we can construct the tableaux  $(P, Q)$  and have proven:

**Theorem 14.** For  $\sigma \in S_n$ , if  $\sigma \leftrightarrow (P, Q)$  then  $\sigma^{-1} \leftrightarrow (Q, P)$ .

