## Huffman Code

Begin with a text file with the following frequencies

| letter | A | B | C | D | E | F | G |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| frequency | 2 | 4 | 6 | 10 | 13 | 13 | 16 |

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| frequency | 2 | 4 | 6 | 10 | 13 | 13 | 16 |
| code length | 5 | 5 | 4 | 3 | 2 | 2 | 2 |

average bits per letter $=(5 \cdot 2+5 \cdot 4+4 \cdot 6+3 \cdot 10$

$$
+2 \cdot 13+2 \cdot 13+2 \cdot 16) / 64=\frac{168}{64}=2.625
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- 

$$
\begin{aligned}
& \text { Entropy }=\frac{2}{64} \log _{2}(32)+\frac{4}{64} \log _{2}(16)+\frac{6}{64} \log _{2}\left(\frac{64}{6}\right) \\
& +\frac{10}{64} \log _{2}\left(\frac{64}{10}\right)+2 \times \frac{13}{64} \log _{2}\left(\frac{64}{13}\right)+\frac{1}{4} \log _{2}(4) \approx 2.579
\end{aligned}
$$

## Tree from heights

Note that given probabilities $p_{A}, p_{B}, \ldots, p_{Z}$, if we set

$$
h_{\alpha}=\left\lceil\log _{2}\left(\frac{1}{p_{\alpha}}\right)\right\rceil
$$

then since we know from Theorem 4 that $\sum_{\alpha=A}^{Z} h_{\alpha} \leq 1$ then by Theorem 1 these values must correspond to heights of a (possibly incomplete) binary tree.

By the same proof as in theorem 4, this code will also have an expected code length less than or equal to $H+1$.

## Tree from heights

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| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| frequency | 2 | 4 | 6 | 10 | 13 | 13 | 16 |

The goal is to encode each letter in such a way that minimizes the average number of bits used to store the file.

Tree from Heights

$$
\begin{array}{c|ccccccc}
\alpha & \text { A } & \text { B } & \text { C } & \text { D } & \text { E } & \text { F } & \text { G } \\
p_{\alpha} & \frac{2}{64} & \frac{4}{64} & \frac{6}{64} & \frac{10}{64} & \frac{13}{64} & \frac{13}{64} & \frac{16}{64} \\
\left\lceil\log _{2}\left(\frac{1}{p_{\alpha}}\right)\right\rceil & 5 & 4 & 4 & 3 & 3 & 3 & 2
\end{array}
$$

## Tree from heights

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| letter | A | B | C | D | E | F | G |
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| frequency | 2 | 4 | 6 | 10 | 13 | 13 | 16 |
| code length | 4 | 4 | 3 | 3 | 3 | 2 | 2 |

average bits per letter $=(4 \cdot 2+4 \cdot 4+3 \cdot 6+3 \cdot 10$

$$
+3 \cdot 13+2 \cdot 13+2 \cdot 16) / 64=\frac{169}{64} \approx 2.641
$$

- 

$$
\begin{aligned}
& \text { Entropy }=\frac{2}{64} \log _{2}(32)+\frac{4}{64} \log _{2}(16)+\frac{6}{64} \log _{2}\left(\frac{64}{6}\right) \\
& +\frac{10}{64} \log _{2}\left(\frac{64}{10}\right)+2 \times \frac{13}{64} \log _{2}\left(\frac{64}{13}\right)+\frac{1}{4} \log _{2}(4) \approx 2.579
\end{aligned}
$$

## Experiment:

Random text consisting of taken from NYTimes consisting of 96,558 alphabetic characters (punctuation and spacing stripped from file).

| A | B | C | D | E | F | G |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7964 | 1466 | 3172 | 3897 | 11547 | 2023 | 1918 |
| H | I | J | K | L | M | N |
| 4626 | 7411 | 292 | 647 | 3955 | 2417 | 7007 |
| O | P | Q | R | S | T | U |
| 7423 | 1966 | 108 | 6113 | 6547 | 8947 | 2715 |
| V | W | X | Y | Z |  |  |
| 1047 | 1565 | 139 | 1532 | 114 |  |  |

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Calculate entropy of this file to be approximately 4.1727 .

Using this text file with 96,558 characters and entropy 4.1727. Using three UNIX file compression programs zip, compress and gzip. I wanted to see how close to the theoretical minimum that I could get.

- compress:
file length $=45,122$ bytes or 360,976 bits. The average number of bits per character is approximately 3.7384.
- gzip:
file length $=39,584$ bytes or 316,672 bits. The average number of bits per character is approximately 3.2796 .
- zip:
file length $=39,706$ bytes or 317,648 bits. The average number of bits per character is approximately 3.2897.
- Wait!? How is it possible? You got better than the theoretical minimum? Oops! Read the instructions, and notice that they are encoding 32 bits at a time (not 8 bits).

Using this text file with $4 \times 96,558$ characters and entropy 4.1727 . Using three UNIX file compression programs zip, compress and gzip. I wanted to see how close to the theoretical minimum that I could get.

- compress:
file length $=62,159$ bytes or 497,272 bits. The average number of bits per character is approximately 5.15 .
- gzip:
file length $=57,404$ bytes or 459,232 bits. The average number of bits per character is approximately 4.76.
- zip:
file length $=57,526$ bytes or 317,648 bits. The average number of bits per character is approximately 4.77.
- Thats better. These values are close (but larger than) the theoretical minimum.

