

$$H(X|Y) \leq H(X)$$

$$H(X, Y) = H(X) + H(Y|X)$$

$$H(X, Y) \leq H(X) + H(Y)$$

$$H(X) \leq \log_2 k \quad \text{if } X \\ \text{takes on } k \\ \text{distinct values}$$

Proof:

$$\begin{aligned} H(X) &= \sum_a P[X=a] \log_2 \left(\frac{1}{P[X=a]} \right) \\ &\leq \log_2 \left(\sum_a P[X=a] \frac{1}{P[X=a]} \right) = \log_2(k) \end{aligned}$$

Unicity Distance for Caesar

Assume that we have just intercepted N letters of ciphertext that was encrypted using a Caesar shift. How large does N have to be (on average) in order to uniquely determine the shift? Assume that the entropy of the English language is 3.2 bits.

We begin with the identity:

$$H(K) = H(C) - H(M)$$

$$H(C) \leq \log_2 26 \approx 4.7$$

$$H(M) < 4.7$$

$$1.78 <$$

Assuming that each of the 26 keys is equally likely, we have

$$H(K) = \log_2 26 \approx 4.7$$

Assuming that each of the 26^N ciphertexts is equally likely, we have

$$H(C) = \log_2 26^N = N \log_2 26 \approx 4.7N$$

Therefore,

$$4.7 = 4.7N - 3.2N \Rightarrow N \approx 3.13$$

Monoalphabetic Substitution

Assume that we have intercepted N letters of a ciphertext message that was encoded using a Monoalphabetic substitution and that the entropy of english is 2 bits. *← I have a larger N for monoalphabetic per letter*

Length of text	5	10	15	20	30	40	50
# of distinct letters	4	8	11	12	14	16	18

For instance, a typical english sample of 30 letters contains about 14 different letters. Thus the key for a Monoalphabetic substitution only permutes 14 letters. Therefore the number of keys is

guess $N \approx 30$

$$26 \times 25 \times \dots \times 13 = (26-14+1)$$

$$(26-1+1)(26-2+1)(26-3+1) \dots (26-14+1)$$

and not $26!$.

Assuming that each key is equally likely, we have

$$H(K) = \log_2(26 \times 25 \times \dots \times 13) \approx 59.54$$

Assuming that each of the 26^N ciphertexts is equally likely, we have

$$H(C) = \log_2 26^N = N \log_2 26 \approx 4.7N$$

Therefore,

$H(M) = 2N$

$$59.54 = 4.7N - 2N \Rightarrow N \approx 22.05$$

With $N \approx 22$ then the # distinct letters is not 14 but closer to 12 or 13

Maybe # of keys =

$$26 \cdot 25 \cdot 24 \cdot \dots \cdot 15$$

OR

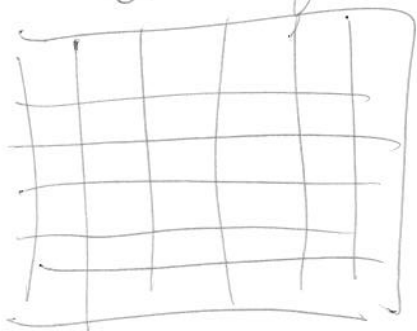
$$26 \cdot 25 \cdot 24 \cdot \dots \cdot 14$$

Try recalculating N for
"this" estimation of $H(K)$

$$H(K) = \log_2(26 \cdot 25 \cdot \dots \cdot 15)$$

ADFGVX - encipherment scheme

Key 1 filling of
6x6 square



Key 2 permutation
of 18

$$\begin{aligned} H(K) &= \log_2 \left(18! \cdot \overset{(36-1+1)}{36 \cdot 35 \cdot 34 \cdot \dots \cdot 11} \right) \\ &= \log_2(18!) + \log_2\left(\frac{36!}{10!}\right) \\ &= 52.5 + 116.3 \end{aligned}$$

$$H(C) = \log_2 6^N = N(1 + \log_2 3)$$

$$H(M) = 1.2(N/2)$$

$$N = \frac{52.5 + 116.3}{(1 + \log_2 3) - 1.2/2}$$