

$a$  &  $\beta$  ( $:= a^{s_B} \pmod{p}$ ) are public for Bob  
 primitive root.  $s_B$  is secret  
 $s_A$  is Alice's secret  
 $a^{s_A} \pmod{p}$  was share  $\beta^{s_A} = (a^{s_A})^{s_B} \pmod{p}$   
**ElGamal Public Key System**

To send a message  $X$  to **Bob** using his public key  $\beta$ , **Alice** chooses at random a secret number  $S_A$  in the interval  $\{1, \dots, p-1\}$ , and sends the pair

"Z" = ciphertext

$(Y, Z)$

where

$Y := a^{s_A} \pmod{p}$ , and  $Z := X \beta^{s_A} \pmod{p}$

Bob knows how to use  $Y$  to decrypt "Z"

↑  
Special key that only Bob can use

↑  
Message

**Bob** can then get  $X$  back using his secret exponent  $S_B$ :

$$X \equiv Z(Y^{s_B})^{-1} \pmod{p}.$$

In this, we can consider that  $Y$  is used to "encode"  $S_A$ .

Note  $Y^{s_B} = (a^{s_A})^{s_B} = (a^{s_B})^{s_A} \equiv \beta^{s_A} \pmod{p}$

which Bob can calculate using his secret key

If Bob knows  $\beta^{s_A}$  then

$$Z \cdot (\beta^{s_A})^{-1} \equiv X \beta^{s_A} \cdot (\beta^{s_A})^{-1} \equiv X \pmod{p}$$

$$a=2 \quad p=17$$

|       |   |   |   |    |    |    |   |   |   |    |    |    |    |    |    |    |
|-------|---|---|---|----|----|----|---|---|---|----|----|----|----|----|----|----|
| $n$   | 1 | 2 | 3 | 4  | 5  | 6  | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| $2^n$ | 2 | 4 | 8 | 16 | 15 | 13 | 9 | 1 |   |    |    |    |    |    |    |    |

$\text{mod } 17$

2 is not a primitive root

|       |   |   |    |    |   |    |    |    |    |    |    |    |    |    |    |    |
|-------|---|---|----|----|---|----|----|----|----|----|----|----|----|----|----|----|
| $n$   | 1 | 2 | 3  | 4  | 5 | 6  | 7  | 8  | 9  | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| $3^n$ | 3 | 9 | 10 | 13 | 5 | 15 | 11 | 16 | 14 | 8  | 7  | 4  | 12 | 2  | 6  | 1  |

3 is a primitive root

Bob has public key  $\beta = 8 \equiv 3^{s_B}$   
 $\Rightarrow s_B = 10$

Alice sends  $(Y, Z) = (12, 10)$   $12 \equiv 3^{s_A}$

Bob calculates  $12 \cdot (10^{s_B})^{-1} \equiv 12 \cdot 9 \equiv 6 \pmod{17}$

$$\begin{aligned} 12 \cdot ((3^3)^{10})^{-1} &\equiv 12 \cdot (3^{30})^{-1} \equiv 12 \cdot 3^{-14} \equiv 12 \cdot 3^2 \\ &\equiv 3^{13} \cdot 3^2 \equiv 3^{15} \end{aligned}$$

Bob's secret key is 10

Alice's secret key is 13

$$8^x \equiv (3^{10})^x \equiv 3^{10x}$$

# Baby step/Giant step method

Goal: Solve  $a^x \equiv b \pmod{n}$ .

Idea: Find  $a^i \equiv ba^{-j} \pmod{n}$  by searching through a small enough space of possible  $i$  and  $j$ .

Idea make a table of  $a^i$  &  $ba^{-j}$   
if I find  $i$  &  $j$  s.t.  $a^i \equiv ba^{-j} \pmod{n}$   
Fix  $m = \lceil \sqrt{(\phi(n))} \rceil$  then find  $c \equiv a^{-m} \pmod{n}$ . Then  $a^{i+j} \equiv b \pmod{n}$

Next calculate a table of  $a^i \pmod{n}$  for  $0 \leq i < m$  and then calculate  $bc^j \pmod{n}$  for  $0 \leq j < m$  until you find one of these values in the table.

Solution: When we find  $a^i \equiv bc^j \pmod{n}$  then we have  $a^{i+mj} \equiv a^i c^{-j} \equiv b \pmod{n}$ .

Example:  $p = 53$  and  $a = 3$ . We wish to solve

$$3^x \equiv 41 \pmod{53}.$$

$\not\vdash a^{-j^m}$  for  $0 \leq j \leq m$

- $m = \lceil \sqrt{\phi(53)} \rceil = 8$  and  $3^{-8} \equiv 24 \pmod{53}$ .

- Now  $41 \cdot 24^i \pmod{53}$ .

| i | $3^i \pmod{53}$               | i | $41 \cdot 24^i \pmod{53}$                                       |
|---|-------------------------------|---|---|
| 0 | 1                             | 0 | 41  |
| 1 | 3                             | 1 | 30  |
| 2 | 9                             | 2 | $\textcircled{31} \equiv 41 \cdot 24^2 \equiv 41 \cdot 3^{-16}$ |
| 3 | 27                            | 3 | 2   |
| 4 | 28                            | 4 | 48  |
| 5 | $\textcircled{31} \equiv 3^5$ | 5 | 39  |
| 6 | 40                            | 6 | 35  |
| 7 | 14                            | 7 | 45  |

$41 \equiv 3^{5+16} \pmod{53}$   
 $\equiv 3^{21} \pmod{53}$

- Conclusion:  $3^{2 \cdot 8 + 5} \equiv 3^{21} \equiv 41 \pmod{53}$

$$1 \leq a \leq p-1 \quad a^{p-1} \equiv 1 \pmod{p}$$

**Theorem 3** For any prime  $p > 2$  and any integer  $a$  not equal to  $0 \pmod{p}$  we have

$$a^{(p-1)/2} = \begin{cases} 1 & \text{if } a \in QR[p] \\ -1 & \text{if } a \notin QR[p] \end{cases}$$

### Proof.

If  $a = x^2$  with  $x \neq 0 \pmod{p}$  then Fermat's theorem gives

$$(x^2)^{(p-1)/2} \equiv a^{(p-1)/2} = x^{p-1} = 1 \pmod{p}$$

Thus the first part of our assertion holds true. To prove the second part, note that the equation

$$x^{p-1} - 1 = 0 \pmod{p}$$

has exactly  $p - 1$  solutions in  $\{1, 2, \dots, p - 1\}$  and for  $p > 2$  we have the factorization

$$x^{p-1} - 1 = (x^{(p-1)/2} - 1)(x^{(p-1)/2} + 1) \equiv 0$$

All  $(p - 1)/2$  elements of  $QR[p]$  satisfy the first factor. Therefore the other  $(p - 1)/2$  solutions must satisfy

$$x^{(p-1)/2} + 1 = 0.$$

## Legendre Symbol

For a prime  $p$

$$\left(\frac{a}{p}\right) = \begin{cases} 1 & \text{if } a \in QR[p] \\ -1 & \text{if } a \notin QR[p] \\ 0 & \text{if } \gcd(a, p) > 1 \end{cases}$$

$L(a, p)$

Then for  $a$  relatively prime to  $p$ , we have

$$\left(\frac{a}{p}\right) = a^{(p-1)/2} \pmod{p} \quad \text{Because of the last theorem.}$$

Hence

$$(ab)^{\frac{p-1}{2}} \equiv \left(\frac{ab}{p}\right) = \left(\frac{a}{p}\right) \left(\frac{b}{p}\right) = a^{\frac{p-1}{2}} b^{\frac{(p-1)}{2}} \pmod{p}$$

**Theorem 4 (Quadratic Reciprocity)** For any two primes  $p$  and  $q$  we have

$$\left(\frac{p}{q}\right) \left(\frac{q}{p}\right) = (-1)^{(p-1)(q-1)/4}$$

~~J(7, 12)~~

$$\cancel{J(7, 12)} \cdot \cancel{J(12 \bmod 7, 7)(+)} -$$

$$J(7, 15) = J(15 \bmod 7, 7) (-1)^{21} = -J(1, 7) = -1$$

by def  $\left(\frac{7}{5}\right) \left(\frac{7}{3}\right) \equiv \left(\frac{2}{5}\right) \left(\frac{1}{3}\right) \equiv (-1) \cdot (1) = -1$

$$J(5, 77) = J(77 \bmod 5, 5)(-1)^{76 \cdot 4/4} = J(2, 5) \equiv J(1, 5)(-1)^{\frac{(25-1)}{8}} = 1 \cdot (-1) = -1$$

$$= \left(\frac{5}{7}\right) \left(\frac{5}{11}\right) = (-1)(+1) = -1$$

$$(\bmod 7)$$

$$1^2 = 1, 2^2 = 4, 3^2 = 2$$

$$1^2 = 1, 2^2 = 4, 3^2 = 9, 4^2 = 5 \pmod{11}$$

## Primality Testing

The Jacobi symbol allows us to test for primality of  $n$  without carrying out its factorization.

If  $n$  is prime then

$$\left(\frac{\alpha}{n}\right) \leftarrow \begin{array}{l} \text{if } n \\ \text{prime} \end{array}$$
$$J(a, n) = a^{(n-1)/2} \pmod{n}$$

Thus if this identity fails to hold for any value of  $a$  in  $[1, n - 1]$  we can certainly conclude that  $n$  is not a prime!

**Theorem 5** *If  $n$  is not a prime then for more than one half the integers in  $\{1, \dots, n - 1\}$  one of the following two tests will fail*

$$J(a, n) = a^{(n-1)/2} \quad \gcd(a, n) = 1$$

## Jacobi Symbol

We start with the Legendre symbol

$$\left(\frac{a}{p}\right) = \begin{cases} 1 & \text{if } a \in QR[p] \\ -1 & \text{if } a \notin QR[p] \end{cases}$$

and for

extension of Legendre symbol to all integers.  
 $n = p_1 p_2 \cdots p_k$

we set

$$J(a, n) = \left(\frac{a}{p_1}\right) \left(\frac{a}{p_2}\right) \cdots \left(\frac{a}{p_k}\right)$$

Jacobi symbol

However, for  $n$  odd, we have

$$J(a, n) = \begin{cases} 1 & \text{if } a = 1 \\ J(a/2, n)(-1)^{(n^2-1)/8} & \text{if } a \text{ is even} \\ J(n \bmod a, a)(-1)^{(n-1)(a-1)/4} & \text{if } a > 1 \text{ and odd} \end{cases}$$

Recursive method for  
computing Jacobi symbol.