Definition $a \equiv b \pmod{m}$ means that m divides a - b or there exists a k such that km = a - b.

 \equiv (mod m) is an equivalence relation since

- $a \equiv a \pmod{m}$ or *m* divides a a. (reflexive)
- if a ≡ b (mod m), then b ≡ a (mod m) since if m divides a - b then it divides b - a. (symmetric)
- if $a \equiv b \pmod{m}$ and $b \equiv c \pmod{m}$, then $a \equiv c \pmod{m}$. (transitive)

It is defined on the integers and so it can easily be shown that for any integer k,

• $a \equiv b \pmod{m}$ if and only if $a + k \equiv b + k \pmod{m}$.

• if
$$a \equiv b \pmod{m}$$
, then $ka \equiv kb \pmod{m}$.

if
$$ka \equiv bk \pmod{m}$$
, then sometimes $a \not\equiv b \pmod{m}$.
e.g. $2 \cdot 3 \equiv 2 \cdot 0 \pmod{6}$, but $3 \not\equiv 0 \pmod{6}$.
e.g. $4 \cdot 5 \equiv 4 \cdot 2 \pmod{12}$, but $5 \not\equiv 2 \pmod{12}$.

When gcd(k, m) = 1, if $ka \equiv kb \pmod{m}$, then $a \equiv b \pmod{m}$

gcd(a, b) = greatest common divisor of a and b = largest divisor of both a and b = if d divides a and b, then d also divides gcd(a, b)

Example: compute *gcd*(963, 657)

 $963 = 1 \cdot 657 + 306$ $657 = 2 \cdot 306 + 45$ $306 = 6 \cdot 45 + 36$ $45 = 1 \cdot 36 + 9$ $36 = 4 \cdot 9$

Conclusion: gcd(963, 657) = 9

$$gcd(963, 657) = 9 = -36 + 45$$

= -(306 - 6 \cdot 45) + 45
= -306 + 7 \cdot 45
= -306 + 7(657 - 2 \cdot 306)
= -15 \cdot 306 + 7 \cdot 657
= -15(963 - 657) + 7 \cdot 657
= -15 \cdot 963 + 22 \cdot 657

In general we can always use these equations to write

$$gcd(a,b) = k \cdot a + \ell \cdot b$$

for some integers k and ℓ .

Example solve $127x \equiv 4 \pmod{963}$

$$963 = 7 \cdot 127 + 74$$

$$127 = 1 \cdot 74 + 53$$

$$74 = 1 \cdot 53 + 21$$

$$53 = 2 \cdot 21 + 11$$

$$21 = 1 \cdot 11 + 10$$

$$11 = 1 \cdot 10 + 1$$

$$1 = 11 - 10 = 11 - (21 - 11) = 2 \cdot 11 - 21$$

= 2(53 - 2 \cdot 21) - 21 = 2 \cdot 53 - 5 \cdot 21 = 2 \cdot 53 - 5(74 - 53)
= 7 \cdot 53 - 5 \cdot 74 = 7(127 - 74) - 5 \cdot 74 = 7 \cdot 127 - 12 \cdot 74
= 7 \cdot 127 - 12(963 - 7 \cdot 127) = 91 \cdot 127 - 12 \cdot 963

Conclusion, because $91 \cdot 127 - 12 \cdot 963 = 1$,

$$91 \cdot 127 \equiv 1 \pmod{963}$$

Therefore if we have

 $127x \equiv 4 \pmod{963}$ $x \equiv 1 \cdot x \equiv 91 \cdot 127x \equiv 91 \cdot 4 \equiv 364 \pmod{963}$

Computational elements that we will use in some new cryptosystems

- Compute $a^k \pmod{m}$ using only squaring operations and multiplication by *a*.
- gcd(a, b) using the Euclidean algorithm
- Find k and ℓ such that

$$ka + \ell b = gcd(a, b)$$

• If gcd(a, m) = 1, then there is a k such that

$$ak \equiv 1 \pmod{m}$$

There is a function called the Euler 'phi' function

 $\phi(n) = \#$ of integers relatively prime (i.e. gcd(k, n) = 1) and are between 1 and n

	integers between 1 and n	
п	which are relatively prime	$\phi(n)$
1	1	1
2	1	1
3	1,2	2
4	1,3	2
5	1,2,3,4	4
6	1,5	2
7	1,2,3,4,5,6	6
8	1,3,5,7	4
9	1,2,4,5,7,8	6
10	1,3,7,9	4
11	1,2,3,4,5,6,7,8,9,10	10
12	1,5,7,11	4
14	1,3,5,9,11,13	6
	•	

Let [a, b] represent the interval of integers $\{a, a + 1, \dots, b - 1, b\}$. Notice that

$$\phi(p) = \#$$
 of integers in [1,p] that have common factor with p
= $\#$ of integers [1,p)
= $p - 1$

Also,

$$\begin{split} \phi(p^k) &= p^k - \# \text{ of integers in } [1,p^k] \text{ divisible by } p \\ &= p^k - \# \text{ of } r \cdot p \text{ where } 1 \leq r \leq p^{k-1} \\ &= p^k - p^{k-1} \end{split}$$

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Say that p does not divide n. Then let h be the number of integers in [1, n] that have a common factor with n.

 $\phi(p^k n) = np^k - \# \text{ of integers in } [1, np^k] \text{ with a common}$ factor with *n* or *p* $= np^k - \# \text{ in } [1, np^k] \text{ with a common factor with } n$ $- \# \text{ in } [1, np^k] \text{ with a common factor with } p$ $+ \# \text{ in } [1, np^k] \text{ with a factor with both } n \text{ and } p$ $= np^k - hp^k - np^{k-1} + hp^{k-1}$ $= (n-h)(p^k - p^{k-1}) = \phi(n)(p^k - p^{k-1})$

if $n = p_1^{a_1} p_2^{a_2} \cdots p_k^{a_k}$ where p_i are all distinct primes, then

$$\phi(n) = (p_1^{a_1} - p_1^{a_1-1})(p_2^{a_2} - p_2^{a_2-1}) \cdots (p_k^{a_k} - p_k^{a_k-k})$$