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Hecke group algebras as degenerate affine Hecke algebras

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AMS Meeting, Claremont, May 4th, 2008

arXiv:0711.1561v1 [math.RT] arXiv:0804.3781v1 [math.RT]

Goals

- Advertise Hecke group algebras:
 - Nice structure and representation theory
 - Connections with NCSF, parking functions, ...
 - Connections with 0-Hecke and affine Hecke algebras
- Where does this structure come from?
- Is it useful?

Coxeter groups



Example (Type A_n : symmetric group \mathfrak{S}_{n+1}) Generators: $(s_i)_{i=1,...,n}$: elementary transpositions

Relations:

$$s_i^2 = 1$$

$$s_i s_j = s_j s_i$$

$$s_i s_{i+1} s_i = s_{i+1} s_i s_{i+1}$$

for all $1 \le i \le n$, for all |i - j| > 1, for all $1 \le i \le n - 1$.

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Coxeter groups

Definition (Coxeter group W) Generators : $(s_i)_{i \in S}$ (simple reflections) Relations: $s_i^2 = 1$ and $\underbrace{s_i s_j \cdots}_{m_{i,j}} = \underbrace{s_j s_i \cdots}_{m_{i,j}}$, for $i \neq j$ Group algebra: $\mathbb{C}[W]$

Example (Type A_n : symmetric group \mathfrak{S}_{n+1}) Generators: $(s_i)_{i=1,...,n}$: elementary transpositions Relations:

$$egin{aligned} &s_i^2 = 1 & ext{for all } 1 \leq i \leq n, \ &s_is_j = s_js_i & ext{for all } |i-j| > 1, \ &s_is_{i+1}s_i = s_{i+1}s_is_{i+1} & ext{for all } 1 \leq i \leq n-1 \end{aligned}$$

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0-(Iwahori)-Hecke algebras



Example (Type A_n)

Generators: $(\pi_i)_{i=1,...,n}$: Relations:

$$\pi_i^2 = \pi_i$$
$$\pi_i \pi_j = \pi_j \pi_i$$
$$\pi_i \pi_{i+1} \pi_i = \pi_{i+1} \pi_i \pi_{i+1}$$

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0-Hecke algebras and bubble sort

Example (Right regular representation for type A)



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Hecke algebras

Take q_1 and q_2 parameters, and set $q := -\frac{q_1}{q_2}$.

Definition (Hecke algebra $H(W)(q_1, q_2)$) Generators : $(T_i)_{i \in S}$ Relations: $(T_i - q_1)(T_i - q_2) = 0$ and $\underbrace{T_i T_j \cdots}_{m_{i,j}} = \underbrace{T_j T_i \cdots}_{m_{i,j}}$, for $i \neq j$ Basis: $(T_w)_{w \in W}$

- At q = 1: group algebra $\mathbb{C}[W]$
- At q = 0: 0-Hecke algebra H(W)(0)
- At q not 0 nor a root of unity: isomorphic to $\mathbb{C}[W]$

Realization of T_i as operator in $End(\mathbb{C}W)$:

$$T_i := (q_1 + q_2)\pi_i - q_1s_i$$

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Hecke group algebras

A silly idea during a brainstorm (Thibon, Novelli, H., T., 2003)

Definition (Hecke group algebra HW of a Coxeter group W)

Glue $\mathbb{C}[W]$ and H(W)(0) on their right regular representations:

$$HW := \langle (\pi_i, s_i)_{i \in S} \rangle \quad \subset \quad \mathsf{End}(\mathbb{C}W)$$

- Any interesting structure?
- Contains all Hecke algebras by construction
- Type A: dimension and dimension of the radical in the Sloane!
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The Hecke group algebra of rank 1

$$W := \{1, s\}$$
 $\mathbb{C}W := \mathbb{C}.1 \oplus \mathbb{C}.s$

Basis $\begin{cases} id = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad s = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \pi = \begin{pmatrix} 0 & 0 \\ 1 & 1 \end{pmatrix} \end{cases}$ Relations $s\pi = \pi, \quad \pi s = \overline{\pi}, \quad \overline{\pi} + \pi = 1 + s$

Dimension 1 simple and projective module

$$(1-s).id = (1-s),$$
 $(1-s).s = -(1-s),$ $(1-s).\pi = 0$

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Structure theory of Hecke group algebras (I)

Definition

Left / right descent sets:

 $D_R(w) := \{i \in S \mid w = w's_i\}$ $D_L(w) := \{i \in S \mid w = s_iw'\}$

 $v \in \mathbb{C}W$ *i-left antisymmetric* if $s_i v = -v$

Theorem (H.,T., 2005)

- Proof: simultaneous triangularity of basis and linear relations
- Straightening relations? Presentation?

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Theorem (H., T., 2005)

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- Straightening relations? Presentation?

• Modules P_I for $I \subset S$:

$$P_I := \{ v \in \mathbb{C}W \mid s_i v = -s_i, \forall i \in I \}$$

- Boolean lattice: $I \subset J \Longrightarrow P_J \subset P_I$
- Combinatorics of descent classes: dim $P_I = |_S^I W|$
- Good basis of ℂW: compatible with restriction on each P_I:

• Par ex:
$$\left\{ v_w := \sum_{w' \in W_{S \setminus D_r(w)}} (-1)^{l(w')} w' w \mid w \in W \right\}$$

• Or the Kazhdan-Lusztig basis to be fancy

Proposition

Realization of the Hecke group algebra as digraph algebra: Basis: $\{e_{w,w'} \mid D_L(w) \subset D_L(w')\}, where e_{w,w'}(v_{w''})\delta_{w'',w}v_{w'}\}$

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Proposition

 $\begin{array}{ll} \mbox{Realization of the Hecke group algebra as digraph algebra:} \\ \mbox{Basis:} & \{e_{w,w'} \mid \mathsf{D}_L(w) \subset \mathsf{D}_L(w')\}, & \mbox{where } e_{w,w'}(v_{w''})\delta_{w'',w}v_{w'} \end{array}$

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Representation theory of Hecke group algebras

Theorem (H., T., 2005)

- e_{w,w}: max. decomposition of id into orthogonal idempotents
- HW Morita equivalent to the poset algebra of boolean lattice
- Projective modules: P_I
- Simple modules: $S_I := P_I / \sum P_J$

Left-antisymmetries on I, left-symmetries on the complement By restriction:

Exactly the Young's ribbon representation of W Exactly the projective modules of H(W)(0)

Question

Representation theory of Hecke group algebras

Theorem (H., T., 2005)

- e_{w,w}: max. decomposition of id into orthogonal idempotents
- HW Morita equivalent to the poset algebra of boolean lattice
- Projective modules: P_I
- Simple modules: $S_I := P_I / \sum P_J$

Left-antisymmetries on I, left-symmetries on the complement By restriction:

Exactly the Young's ribbon representation of W Exactly the projective modules of H(W)(0)

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Recursive sorting algorithm in type A

 π_1, \ldots, π_n can antisort 12345 to 54321 (bubble sort) But not back! (going down the permutohedron)

Definition (Affine action)

Put the permutation on a circle π_0 acts between last and first letter

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Proposition (S.,T.,2007)

- Similar algorithms for types B, C, D
- Existence for all types (including twisted):

Proof

• Type B: $1 < 2 < 3 < 4 < \underline{4} < \underline{3} < \underline{2} < \underline{1}$

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Proposition (S.,T.,2007)

- Similar algorithms for types B, C, D
- Existence for all types (including twisted):

Proof

• Type B: $1 < 2 < 3 < 4 < \underline{4} < \underline{3} < \underline{2} < \underline{1}$

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Type free geometric argument (II)



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Hecke group algebras and principal series representations

- $Y^{\lambda^{\vee}} \in H(W)(q)$ (analog of translations in W)
- $\mathbb{C}[Y]$: commutative algebra, $\mathbb{C}[Y]^{\mathring{W}}$: center of H(W)(q)
- t: character on C[Y]
 - Central specialization of $\mathbb{C}[Y]^W$ on t: Quotient $\mathcal{H}(q, t)$ of H(W)(q) of dimension $|W|^2$
 - Representation
 ρ_t := t ↑^{H(W)(q)}_{C[Y]} of H(W)(q) on CW: Quotient of H(q, t), generically trivial

Theorem (S.,T., 2008)

Take q non zero and non root of unity, and $t : Y^{\lambda^{\vee}} \mapsto q^{-\operatorname{ht}(\lambda^{\vee})}$ Then, $\rho_t(\operatorname{H}(W)(q)) = \operatorname{H} \mathring{W}$ I.e. $\operatorname{H} \mathring{W}$ (non trivial!) quotient of $\mathcal{H}(q, t)$ and of $\operatorname{H}(W)(q)$

- Proof: diagonalization of the action of Y on C W, using alcove walks and the intertwining operators τ_i
- What happens at roots of unity?

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