

Plaintext  $\xrightarrow{\text{mono}}$  Cyphertext  $\xrightarrow{\text{vig}}$  Cyphertext

18 distinct      18 distinct      26 distinct

$$\# \text{ of monoalphabetic keys} = \frac{26 \cdot 25 \cdot \dots \cdot 9}{1} = \frac{26!}{(26-18)!}$$

assuming 18 distinct letters

$$= \frac{26!}{8!}$$

$$\# \text{ Vigenere keys w/6 letters} = 26^6$$

$$\text{total \# of keys} = \frac{26^6 \cdot \frac{26!}{8!}}{26}$$

$$\# \text{ Vigenere first} = 26^6 \text{ keys}$$

$$\# \text{ monoalphabetic keys} = 26!$$

$$\text{Total \# of keys} = \frac{26^6 \cdot 26!}{26}$$

~~by~~ Plaintext  $\xrightarrow{\text{vig}}$  Cyphertext  $\xrightarrow{\text{mono}}$  Cyphertext

18 distinct      26 distinct      26 distinct

plaintext, 18 different chars  
monalph  
cypher text 18 diff chars  
Vigenere  
cypher text with 26 possible

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plaintext 18 different chars  
Vigenere  
cypher text 26 different  
monalph bet 26  
cypher text with 26 possible again

Vernam with key  $P = (P_1, P_2, P_3)$   
 $q = (q_1, q_2, q_3, q_4)$

0	10	15	0	10	15	0	20	15	0	20	15	0	20	15
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$P_1$	$P_2$	$P_3$	$P_1$	$P_2$	$P_3$	$P_1$	$P_2$	$P_3$	$P_1$	$P_2$	$P_3$	$P_1$	$P_2$	$P_3$
7	2	14	5	7	2	14	5	7	2	14	5	7	2	14

$q_1$   $q_2$   $q_3$   $q_4$   $q_1$   $q_2$   $q_3$   $q_4$   $q_1$   $q_2$   $q_3$   $q_4$   $q_1$   $q_2$   $q_3$

7	22	3	5	1	17	14	25	22	2	8	20	7	22	3
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$K_1$

0	13	0	17	19	8	18	19	13	4	21	4	17	17	4
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A	W	A	R	T	I	S	T	N	E	V	E	R	R	E
7	9	3	22	20	25	6	18	9	6	3	24	24	13	7

H I J D W U Z G S J G D Y Y N H

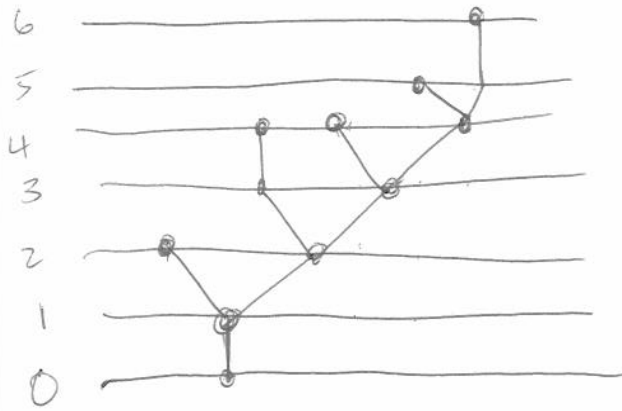
A	0	F	5	K	10	P	15	U	20	Z	25
B	1	G	6	L	11	Q	16	V	21		
C	2	H	7	M	12	R	17	W	22		
D	3	I	8	N	13	S	18	X	23		
E	4	J	9	O	14	T	19	Y	24		

more:  
 given Vernam 3 and length of keys  
 F M C M E E P Q M D A V M J T K N N M

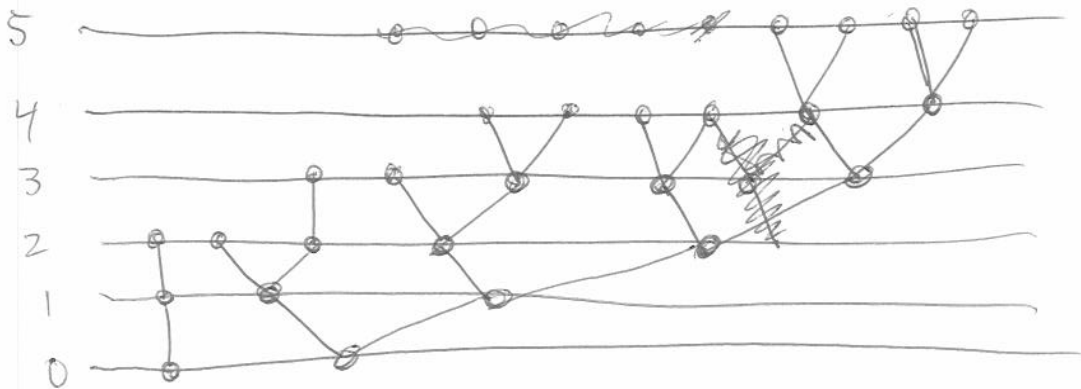
If we have a list of heights

~~At least~~

2 4 4 5 6



2 2 3 3 4 4 4 4 5 5 5 5



If we have a list of heights

$$h_1, h_2, h_3, \dots, h_n$$

if  $\sum \frac{1}{2^{h_i}} \leq 1$  then  
it makes a binary tree

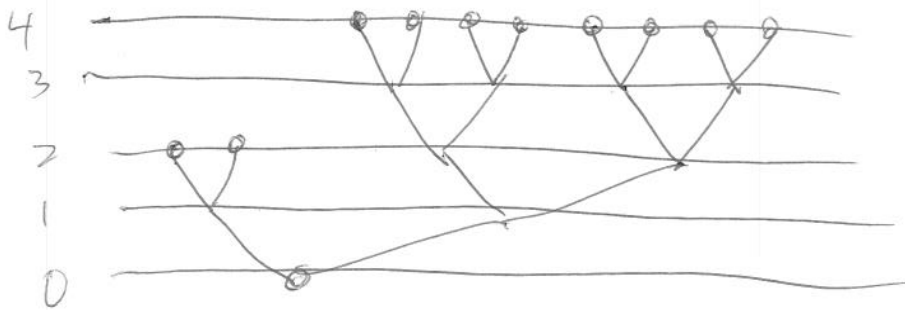
if  $\sum \frac{1}{2^{h_i}} > 1$  then  
it is not a binary tree

if  $\sum \frac{1}{2^{h_i}} = 1$  then  
the binary tree is complete.

$$\frac{1}{2^2} + \frac{1}{2^2} + \frac{1}{2^4} + \frac{1}{2^4} + \frac{1}{2^4} + \frac{1}{2^4} + \frac{1}{2^4} + \frac{1}{2^4}$$

$$\underbrace{\frac{1}{2}}_{\frac{1}{2}} + \frac{1}{2^4} + \frac{1}{2^4} = \frac{1}{2}$$

$$\frac{1}{4} + \frac{1}{4} + \frac{1}{16} + \frac{1}{16} + \frac{1}{16} + \frac{1}{16} + \frac{1}{16} + \frac{1}{16} + \frac{1}{16} + \frac{1}{16} = 1$$



there are 14 possible events in  
 this scenario 7 coins one is lighter  
 OR 7 coins one is heavier

$$H(\text{event}) = \log_2 14 = 1 + \log_2 7 \approx 3.8$$

something  
 = the amount of information  
 you learn about the  
 coin



gives  $\log_2 3$  bits per  
 weighing.  $\log_2 3 \approx 1.7$

